

# Recent advances around Nivat's conjecture

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Combinatorics on Words seminar

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# Nivat's Conjecture

# Pattern Complexity – 1D

$$\mathcal{A} = \{ \text{blue square} \text{ green square} \text{ purple square} \text{ orange square} \}$$

$$w \in \mathcal{A}^{\mathbb{Z}}$$



# Pattern Complexity – 1D

$P_w(n)$  = number of patterns of size  $n$



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$$P_w(1) = 4$$

$$P_w(2) = 11$$

$$P_w(4) = 18$$

## Pattern Complexity – 1D: Periodicity



$$P_w(1) = 2$$

$$P_w(2) = 3$$

$$P_w(3) = 5$$

$$P_w(4) = 5$$

$$P_w(5) = 5$$

$$P_w(6) = 5$$

⋮

## Pattern Complexity – 1D: Periodicity



Property

If  $w$  is periodic, then  $\exists n, P_w(n) \leq n$ .

# Pattern Complexity – 1D

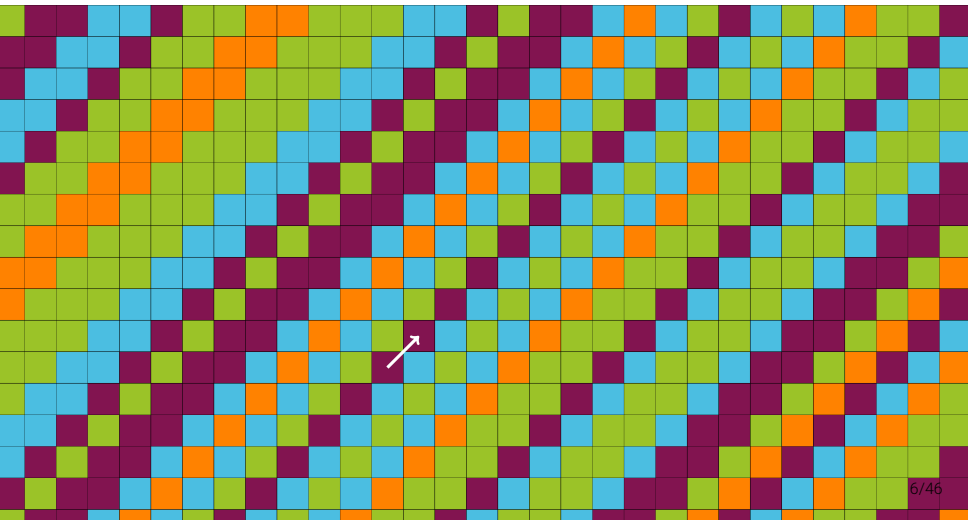
Theorem (Morse & Hedlund, 1938)

$$\forall w \in \mathcal{A}^{\mathbb{Z}},$$
$$\exists n > 0, P_w(n) \leq n \Rightarrow w \text{ periodic}$$

# Periodic Configuration – 2D

$c \in \mathcal{A}^{\mathbb{Z}^2}$  is:

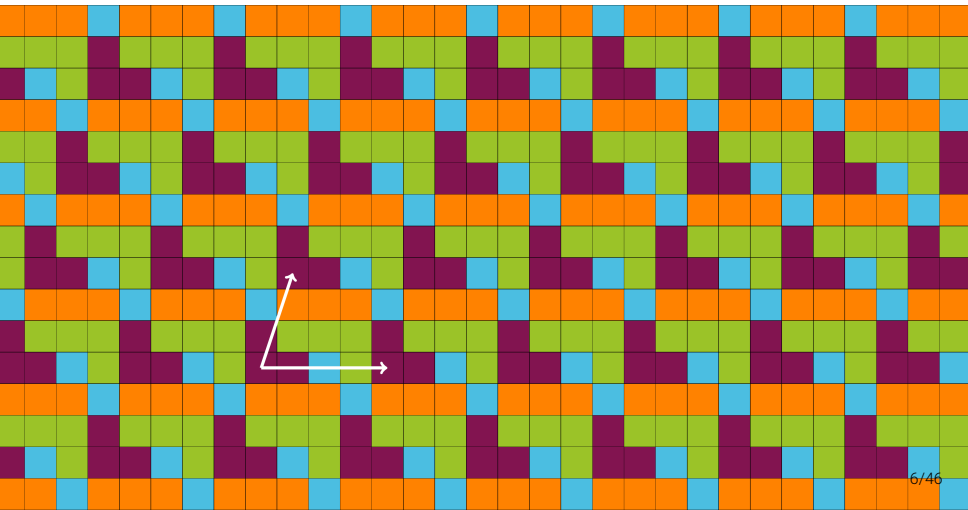
- **1-periodic / weakly periodic:**  $\exists \mathbf{u}, \forall \mathbf{v}, c_{\mathbf{v}-\mathbf{u}} = c_{\mathbf{v}}$



# Periodic Configuration – 2D

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- **1-periodic / weakly periodic:**  $\exists \mathbf{u}, \forall \mathbf{v}, c_{\mathbf{v}-\mathbf{u}} = c_{\mathbf{v}}$
- **2-periodic / strongly periodic:**  $c$  is 1-periodic along  $\mathbf{u}_1, \mathbf{u}_2$   
not colinear

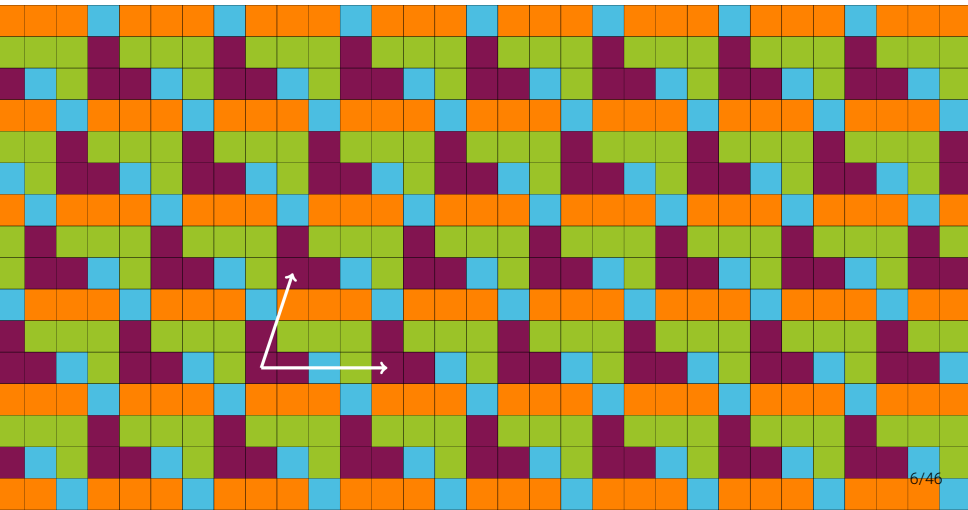




# Periodic Configuration – 2D

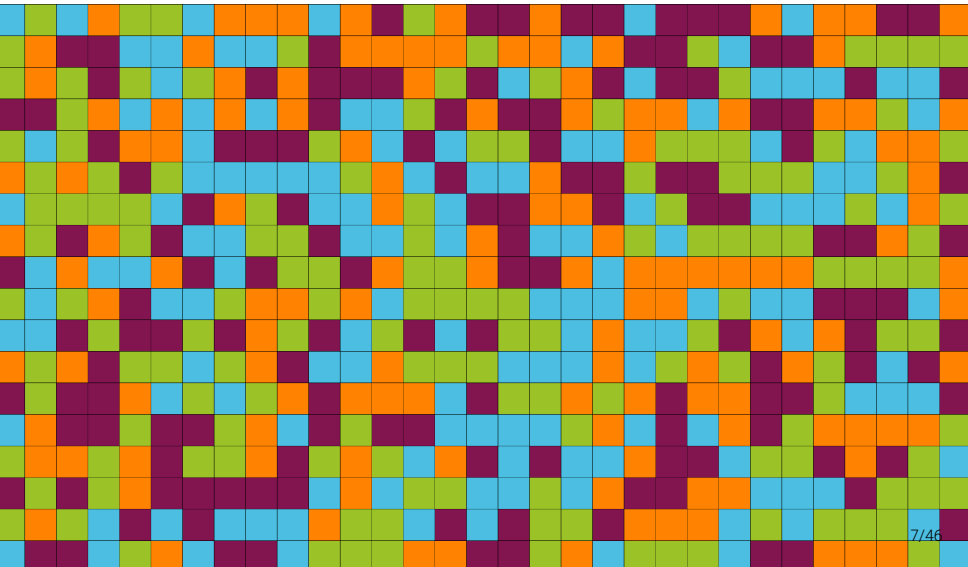
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- **2-periodic / strongly periodic:**  $c$  is 1-periodic along  $\mathbf{u}_1, \mathbf{u}_2$   
not colinear  $\Rightarrow$  Finitely many different translations of  $c$



# Pattern Complexity – 2D

$P_c(m, n)$  = number of rectangular patterns of size  $m \times n$



## 2D: Nivat's conjecture

Conjecture (Nivat, 1997)

$$\forall c \in \mathcal{A}^{\mathbb{Z}^2},$$
$$\exists m, n > 0, P_c(m, n) \leq mn \Rightarrow c \text{ periodic}$$

## 2D: no equivalence

Property

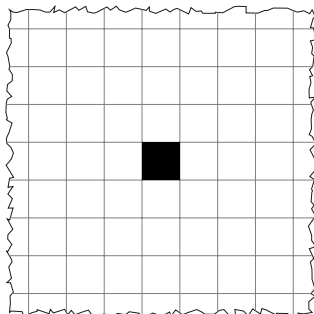
There exists  $c \in \{0, 1\}^{\mathbb{Z}^2}$  periodic s.t.

$$P_c(m, n) = 2^{m+n-1}$$

# Nivat's conjecture is optimal

Theorem

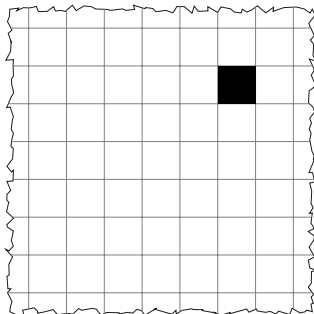
$$\exists c \in \mathcal{A}^{\mathbb{Z}^2},$$
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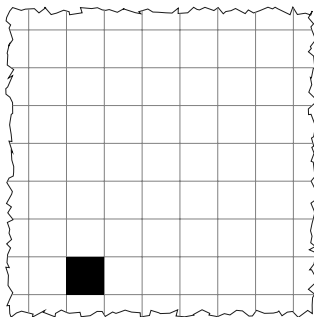
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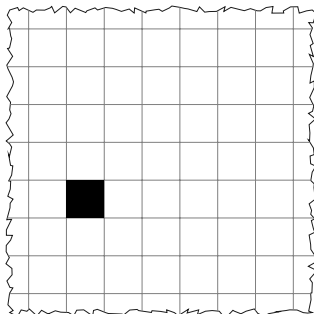
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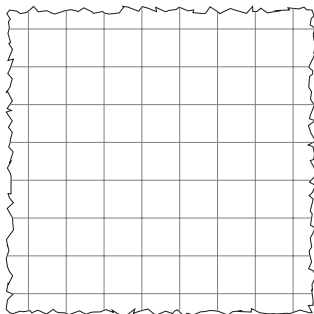




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# Uniform Recurrence

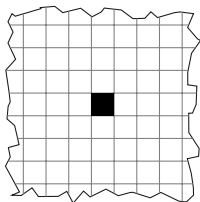
$c$  is **uniformly recurrent** if all patterns of  $c$  appear *everywhere*:

$$\forall p \sqsubset c, \exists n_p, p \text{ appears in all balls of size } n_p$$

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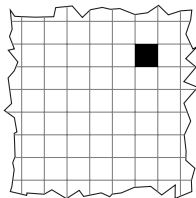


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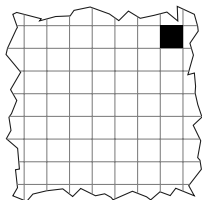


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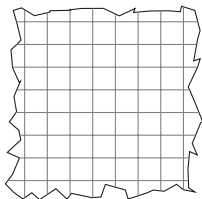


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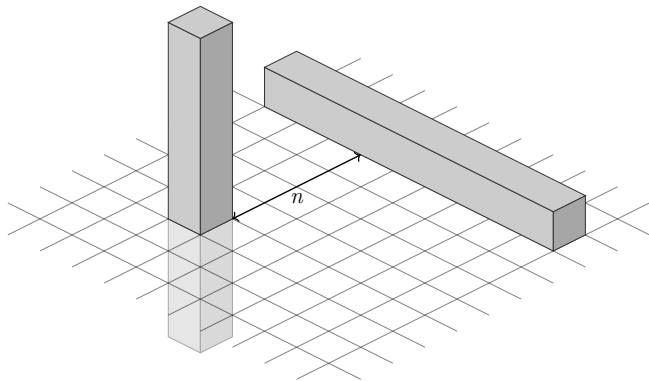
is not uniformly recurrent

## Higher dimension?

The same result does not hold in 3D and above:

$$\exists c \in \mathcal{A}^{\mathbb{Z}^3},$$
$$\exists n > 0, \quad P_c(n, n, n) \leq n^3 \quad \text{and} \quad c \text{ **not** periodic}$$

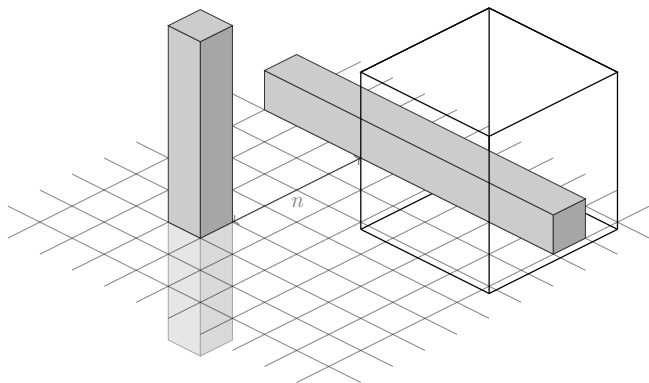
## Higher dimension?



$$P_c(n, n, n) = 2n^2 + 1 < n^3$$

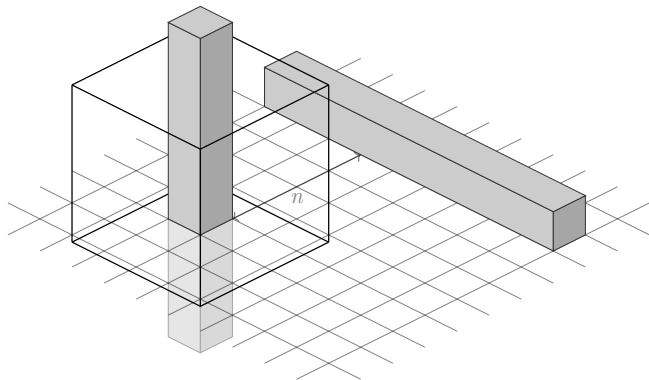


## Higher dimension?



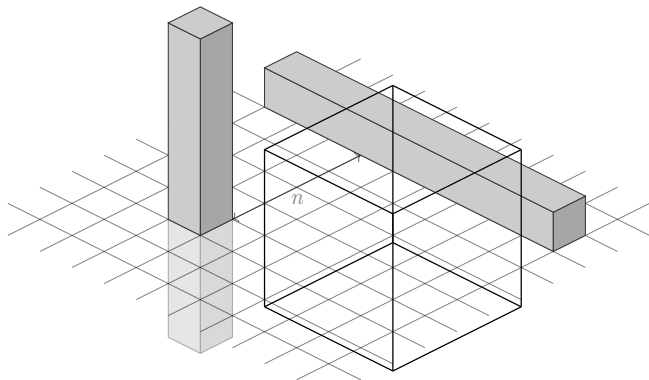
$$P_c(n, n, n) = \mathbf{n^2} + n^2 + 1 < n^3$$

## Higher dimension?



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## Higher dimension?



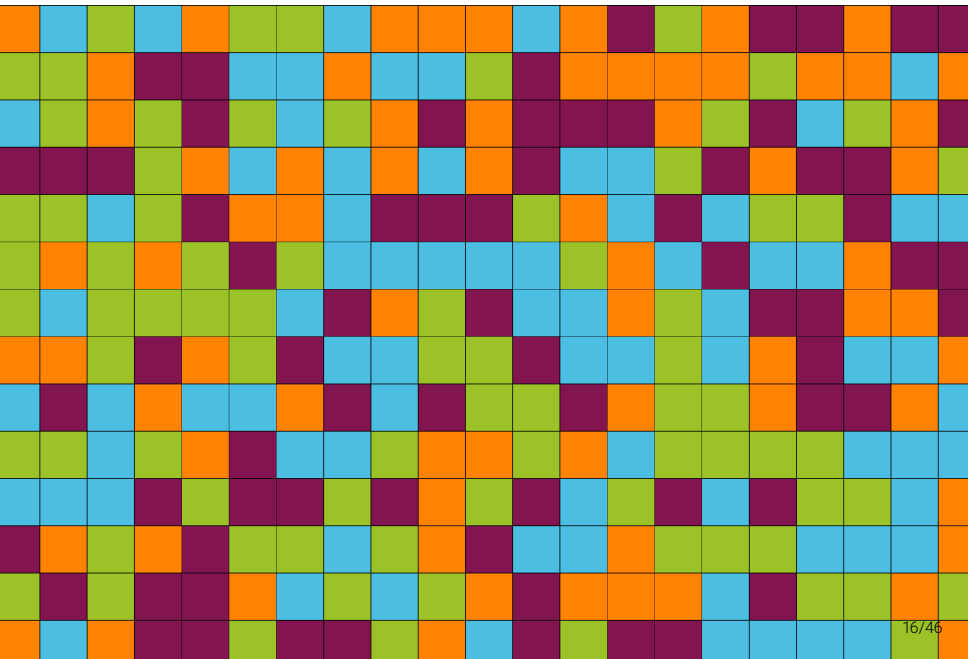
$$P_c(n, n, n) = n^2 + n^2 + \mathbf{1} < n^3$$

## Getting closer to the conjecture

- $P_c(2, n) \leq 2n$  [Sanders & Tijdeman 2002]
- $P_c(3, n) \leq 3n$  [Cyr & Kra, 2016]
  
- $P_c(m, n) \leq \frac{mn}{144}$  [Epifanio, Koskas & Mignosi, 2003]
- $P_c(m, n) \leq \frac{mn}{16}$  [Quas & Zamboni, 2004]
- $P_c(m, n) \leq \frac{mn}{2}$  [Cyr & Kra, 2015]
- $P_c(m, n) \leq \frac{mn}{2} + |\mathcal{A}| - 1$  [Colle & Garibaldi, 2019]

# Algebraic Tools

# Configurations are (Laurent) series





# Configurations are (Laurent) series

-10,5	$c_{-9,5}$	$c_{-8,5}$	$c_{-7,5}$	$c_{-6,5}$	$c_{-5,5}$	$c_{-4,5}$	$c_{-3,5}$	$c_{-2,5}$	$c_{-1,5}$	$c_{0,5}$	$c_{1,5}$	$c_{2,5}$	$c_{3,5}$	$c_{4,5}$	$c_{5,5}$	$c_{6,5}$	$c_{7,5}$	$c_{8,5}$	$c_{9,5}$	$c_{10,5}$
-10,4	$c_{-9,4}$	$c_{-8,4}$	$c_{-7,4}$	$c_{-6,4}$	$c_{-5,4}$	$c_{-4,4}$	$c_{-3,4}$	$c_{-2,4}$	$c_{-1,4}$	$c_{0,4}$	$c_{1,4}$	$c_{2,4}$	$c_{3,4}$	$c_{4,4}$	$c_{5,4}$	$c_{6,4}$	$c_{7,4}$	$c_{8,4}$	$c_{9,4}$	$c_{10,4}$
-10,3	$c_{-9,3}$	$c_{-8,3}$	$c_{-7,3}$	$c_{-6,3}$	$c_{-5,3}$	$c_{-4,3}$	$c_{-3,3}$	$c_{-2,3}$	$c_{-1,3}$	$c_{0,3}$	$c_{1,3}$	$c_{2,3}$	$c_{3,3}$	$c_{4,3}$	$c_{5,3}$	$c_{6,3}$	$c_{7,3}$	$c_{8,3}$	$c_{9,3}$	$c_{10,3}$
-10,2	$c_{-9,2}$	$c_{-8,2}$	$c_{-7,2}$	$c_{-6,2}$	$c_{-5,2}$	$c_{-4,2}$	$c_{-3,2}$	$c_{-2,2}$	$c_{-1,2}$	$c_{0,2}$	$c_{1,2}$	$c_{2,2}$	$c_{3,2}$	$c_{4,2}$	$c_{5,2}$	$c_{6,2}$	$c_{7,2}$	$c_{8,2}$	$c_{9,2}$	$c_{10,2}$
-10,1	$c_{-9,1}$	$c_{-8,1}$	$c_{-7,1}$	$c_{-6,1}$	$c_{-5,1}$	$c_{-4,1}$	$c_{-3,1}$	$c_{-2,1}$	$c_{-1,1}$	$c_{0,1}$	$c_{1,1}$	$c_{2,1}$	$c_{3,1}$	$c_{4,1}$	$c_{5,1}$	$c_{6,1}$	$c_{7,1}$	$c_{8,1}$	$c_{9,1}$	$c_{10,1}$
-10,0	$c_{-9,0}$	$c_{-8,0}$	$c_{-7,0}$	$c_{-6,0}$	$c_{-5,0}$	$c_{-4,0}$	$c_{-3,0}$	$c_{-2,0}$	$c_{-1,0}$	$c_{0,0}$	$c_{1,0}$	$c_{2,0}$	$c_{3,0}$	$c_{4,0}$	$c_{5,0}$	$c_{6,0}$	$c_{7,0}$	$c_{8,0}$	$c_{9,0}$	$c_{10,0}$
-10,-1	$c_{-9,-1}$	$c_{-8,-1}$	$c_{-7,-1}$	$c_{-6,-1}$	$c_{-5,-1}$	$c_{-4,-1}$	$c_{-3,-1}$	$c_{-2,-1}$	$c_{-1,-1}$	$c_{0,-1}$	$c_{1,-1}$	$c_{2,-1}$	$c_{3,-1}$	$c_{4,-1}$	$c_{5,-1}$	$c_{6,-1}$	$c_{7,-1}$	$c_{8,-1}$	$c_{9,-1}$	$c_{10,-1}$
-10,-2	$c_{-9,-2}$	$c_{-8,-2}$	$c_{-7,-2}$	$c_{-6,-2}$	$c_{-5,-2}$	$c_{-4,-2}$	$c_{-3,-2}$	$c_{-2,-2}$	$c_{-1,-2}$	$c_{0,-2}$	$c_{1,-2}$	$c_{2,-2}$	$c_{3,-2}$	$c_{4,-2}$	$c_{5,-2}$	$c_{6,-2}$	$c_{7,-2}$	$c_{8,-2}$	$c_{9,-2}$	$c_{10,-2}$
-10,-3	$c_{-9,-3}$	$c_{-8,-3}$	$c_{-7,-3}$	$c_{-6,-3}$	$c_{-5,-3}$	$c_{-4,-3}$	$c_{-3,-3}$	$c_{-2,-3}$	$c_{-1,-3}$	$c_{0,-3}$	$c_{1,-3}$	$c_{2,-3}$	$c_{3,-3}$	$c_{4,-3}$	$c_{5,-3}$	$c_{6,-3}$	$c_{7,-3}$	$c_{8,-3}$	$c_{9,-3}$	$c_{10,-3}$
-10,-4	$c_{-9,-4}$	$c_{-8,-4}$	$c_{-7,-4}$	$c_{-6,-4}$	$c_{-5,-4}$	$c_{-4,-4}$	$c_{-3,-4}$	$c_{-2,-4}$	$c_{-1,-4}$	$c_{0,-4}$	$c_{1,-4}$	$c_{2,-4}$	$c_{3,-4}$	$c_{4,-4}$	$c_{5,-4}$	$c_{6,-4}$	$c_{7,-4}$	$c_{8,-4}$	$c_{9,-4}$	$c_{10,-4}$

$$c(X, Y) = \sum_{i,j=-\infty}^{\infty} c_{i,j} X^i Y^j$$



## Operations: Sum

$$c + d = \sum_{i,j=-\infty}^{\infty} (c_{i,j} + d_{i,j}) X^i Y^j$$

Formal sum  $\leftrightarrow$  Sum of configurations

## Operations: Multiplication

$$X^a Y^b c = \sum_{i,j=-\infty}^{\infty} c_{i,j} X^{i+a} Y^{j+b} = \sum_{i,j=-\infty}^{\infty} c_{i-a,j-b} X^i Y^j$$

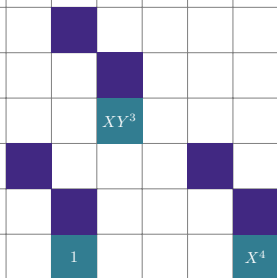
Multiplication by  $X^a Y^b \leftrightarrow$  Translation of vector  $(a, b)$

(Parenthesis: see your polynomials like you never did)

	$Y^2$	
		$XY$
$(0,0)$		$X$

$$X + XY + Y^2$$

(Parenthesis: see your polynomials like you never did)



$$(X + XY + Y^2)(1 + X^4 + XY^3)$$

# Expressing periodicity

$c$  periodic of period  $(a, b)$

$\Leftrightarrow$

$$c = X^a Y^b c$$

$\Leftrightarrow$

$$(X^a Y^b - 1)c = 0$$

$\Leftrightarrow$

$$(X^a Y^b - 1) \in \text{Ann}(c)$$

# Polynomial Ideal

$$\text{Ann}(c) = \{p \mid pc = 0\} \subset \mathbb{R}[X^\pm, Y^\pm]$$

$$c \text{ periodic} \Leftrightarrow \exists a, b \in \mathbb{Z} \setminus \{0\}, (X^a Y^b - 1) \in \text{Ann}(c)$$

# Polynomial Ideal

$$\text{Ann}(c) = \{p \mid pc = 0\} \subset \mathbb{R}[X^{\pm}, Y^{\pm}]$$

$$c \text{ periodic} \Leftrightarrow \exists a, b \in \mathbb{Z} \setminus \{0\}, (X^a Y^b - 1) \in \text{Ann}(c)$$

$\text{Ann}(c)$  is a **polynomial ideal**

→ a lot of tools to understand its structure !

# First results

$c$  of low complexity :  $\exists m, n, P_c(m, n) \leq mn$

Theorem (Kari & Szabados, 2015)

$\exists p \neq 0 \in \text{Ann}(c)$



# First results

$c$  of low complexity :  $\exists m, n, P_c(m, n) \leq mn$

Theorem (Kari & Szabados, 2015)

$$\exists p \neq 0 \in \text{Ann}(c)$$

Theorem (Kari & Szabados, 2015)

$$\exists a_1, b_1, a_2, b_2 \dots, a_r, b_r \in \mathbb{Z}$$

$$\left( X^{a_1} Y^{b_1} - 1 \right) \left( X^{a_2} Y^{b_2} - 1 \right) \dots \left( X^{a_r} Y^{b_r} - 1 \right) \in \text{Ann}(c)$$

# Periodic Decomposition

$c$  of low complexity :  $\exists m, n, P_c(m, n) \leq mn$

Theorem (Kari & Szabados, 2015)

$\exists a_1, b_1, a_2, b_2, \dots, a_r, b_r \in \mathbb{Z}$

$$\left( X^{a_1} Y^{b_1} - 1 \right) \left( X^{a_2} Y^{b_2} - 1 \right) \cdots \left( X^{a_r} Y^{b_r} - 1 \right) \in \text{Ann}(c)$$

Theorem (Kari & Szabados, 2015)

There exist periodic  $c_1, \dots, c_r$

$$c = c_1 + \cdots + c_r$$

# Periodic Decomposition: any dimension !

$c$  in dimension  $d$  with a non-trivial annihilator

Theorem (Kari & Szabados, 2015)

$$\exists \mathbf{u}_1, \dots, \mathbf{u}_r \in \mathbb{Z}^d$$

$$(X^{\mathbf{u}_1} - 1)(X^{\mathbf{u}_2} - 1) \cdots (X^{\mathbf{u}_r} - 1) \in \text{Ann}(c)$$

Theorem (Kari & Szabados, 2015)

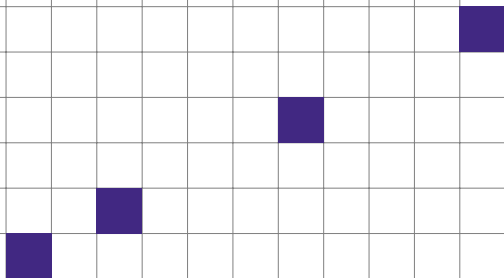
There exist periodic  $c_1, \dots, c_r$

$$c = c_1 + \cdots + c_r$$

# Line Polynomial

$p$  is a **line polynomial** if  $\exists a, b, c, d,$

$$p = X^c Y^d \sum_i p_i X^{ia} Y^{ib}$$



$$1 + X^2 Y + X^6 Y^3 + X^{10} Y^5$$

# Annihilator Ideal Decomposition (2D)

Theorem (Kari & Szabados, 2015)

$c \in \mathcal{A}^{\mathbb{Z}^2}$  w. non-trivial annihilator. Then there are line polynomials  $\phi_1, \dots, \phi_r$  and ideal  $H$  s.t.

$$\text{Ann}(c) = \phi_1 \cdots \phi_r H,$$

$H$  intersection of maximal ideals,  $\langle \phi_1 \cdots \phi_r \rangle$  and  $H$  co-maximal,  $\phi_1, \dots, \phi_r, H$ , unique

# Periodic Decomposition

Theorem (Kari & Szabados, 2015)

If  $c = c_1 + c_2$ ,  $c_1$  and  $c_2$  periodic, and  $P_c(m, n) \leq mn$ , then  $c$  is periodic.

Theorem (Cyr & Kra, 2015)

If  $\exists m, n \in \mathbb{N}$ ,  $P_c(m, n) \leq \frac{mn}{2}$ , then  $c$  is periodic.

# Asymptotic Version

Theorem (Kari & Szabados, 2015)

If  $\exists$  infinitely many  $m, n \in \mathbb{N}$ ,  $P_c(m, n) \leq mn$ , then  $c$  is periodic.

# Uniformly Recurrent Case



# Uniformly recurrent configurations

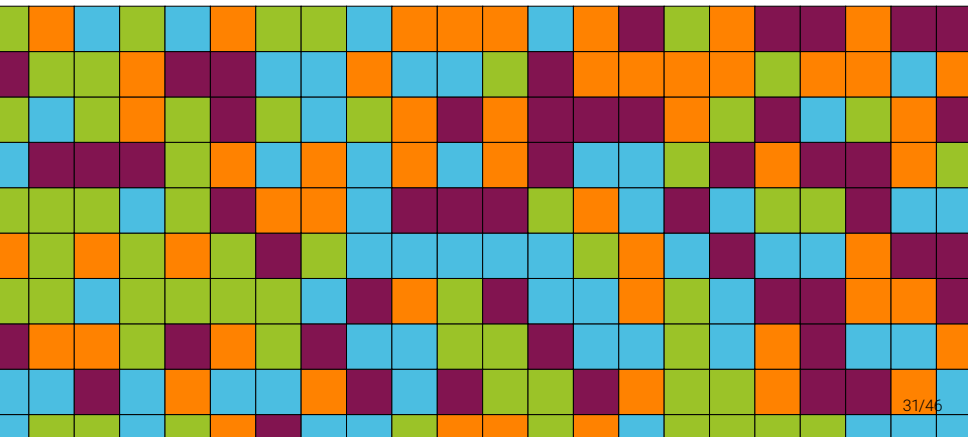
Theorem (Kari & M. 2020)

Nivat's conjecture holds for uniformly recurrent configurations

# Subshifts of Finite Type – Configurations

Finite alphabet:  $\mathcal{A} = \left\{ \begin{array}{c} \text{light blue} \\ \text{light green} \\ \text{dark purple} \\ \text{orange} \end{array} \right\}$

Configuration:  $c \in \mathcal{A}^{\mathbb{Z}^2}$



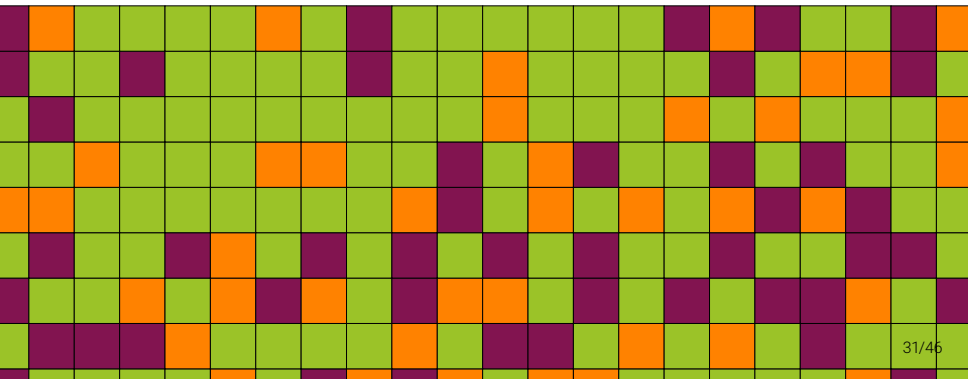
# Subshifts of Finite Type

Finite alphabet:  $\mathcal{A} = \{\text{blue}, \text{green}, \text{purple}, \text{orange}\}$

Set of forbidden patterns:  $F = \{\text{blue}\}$

Subshift :

$$X_F = \{c \in \mathcal{A}^{\mathbb{Z}^2} \mid \forall m \in F, m \text{ does not appear in } c\}$$



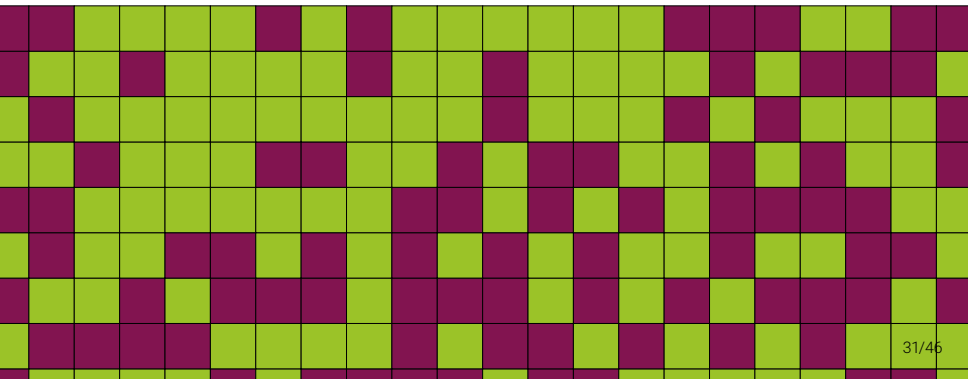
# Subshifts of Finite Type

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Set of forbidden patterns:  $F = \{\text{blue}, \text{orange}\}$

Subshift :

$$X_F = \{c \in \mathcal{A}^{\mathbb{Z}^2} \mid \forall m \in F, m \text{ does not appear in } c\}$$



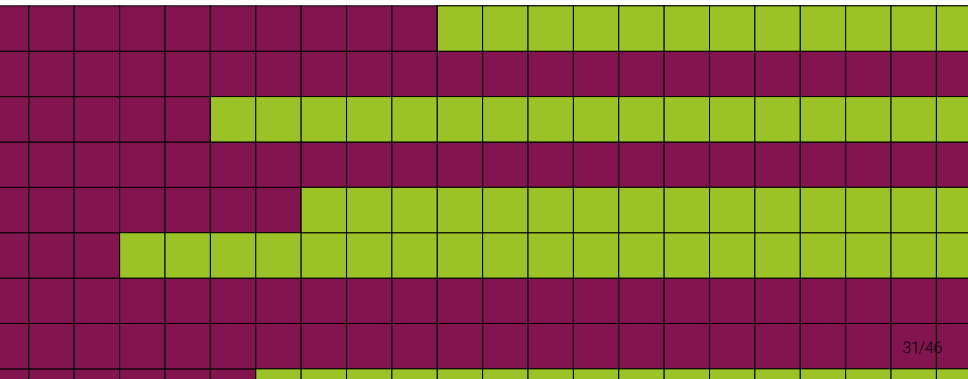
# Subshifts of Finite Type

Finite alphabet:  $\mathcal{A} = \{\text{blue}, \text{green}, \text{purple}, \text{orange}\}$

Set of forbidden patterns:  $F = \{\text{blue orange}, \text{green purple}\}$

Subshift :

$$X_F = \{c \in \mathcal{A}^{\mathbb{Z}^2} \mid \forall m \in F, m \text{ does not appear in } c\}$$



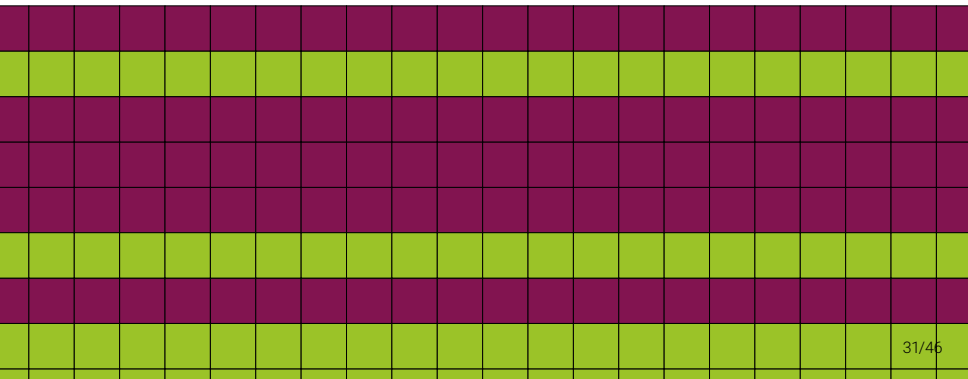
# Subshifts of Finite Type

Finite alphabet:  $\mathcal{A} = \{\text{blue}, \text{green}, \text{purple}, \text{orange}\}$

Set of forbidden patterns:  $F = \{\text{blue orange}, \text{green purple}, \text{purple green}\}$

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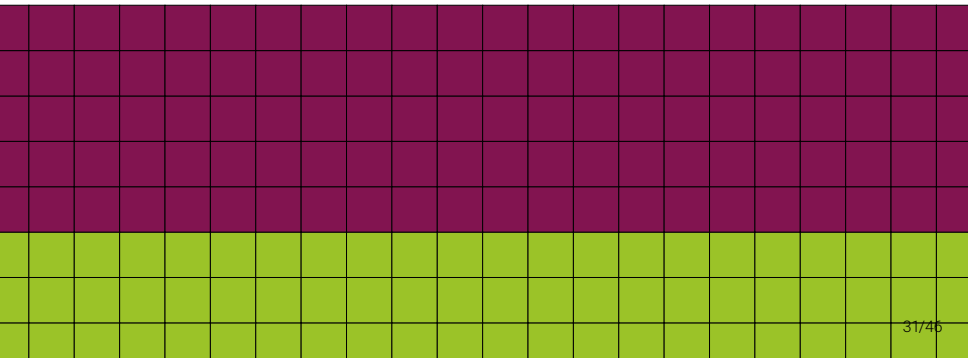
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Subshift :  $X_F$

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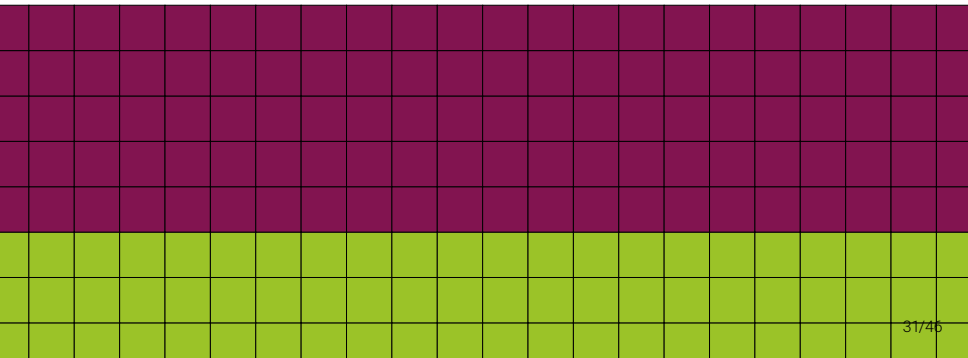
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Subshift of Finite Type (SFT):

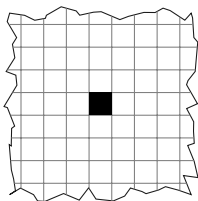
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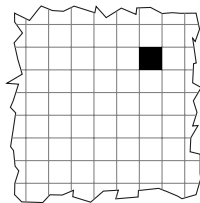
# Orbits

**Orbit** of  $c$ :  $\mathcal{O}(c) = \{\mathbf{u} \cdot c \mid \mathbf{u} \in \mathbb{Z}^2\}$



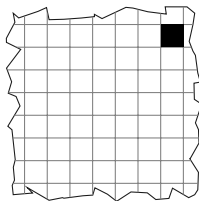
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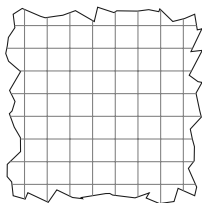
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**Orbit closure** of  $c$ :  $\overline{\mathcal{O}(c)}$ , topological closure of  $\mathcal{O}(c)$  (a subshift)

$c$  uniformly recurrent  $\Leftrightarrow \forall d \in \overline{\mathcal{O}(c)}, \overline{\mathcal{O}(d)} = \overline{\mathcal{O}(c)}$

# Main Theorem

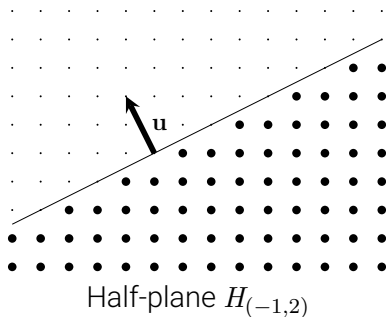
Theorem (Kari & M. 2020)

$c$  such that  $\exists m, n, P_c(m, n) \leq mn,$

$\exists d \in \overline{\mathcal{O}(c)}$  which is periodic

Key: **determinism**

# Determinism



$X$  is  $\mathbf{u}$ -deterministic:

$$\forall c, c' \in X,$$
$$c|_{H_{\mathbf{u}}} = c'|_{H_{\mathbf{u}}} \implies c = c'$$

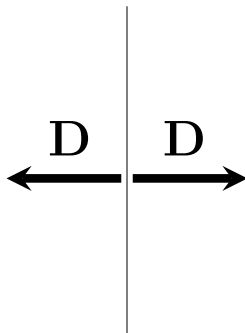
## Why Determinism? (1)

Lemma (Corollary from [Boyle & Lind 1997])

$\overline{\mathcal{O}(c)}$  is deterministic in all directions

$\Rightarrow$

$c$  is (two-)periodic





## Why Determinism? (2)

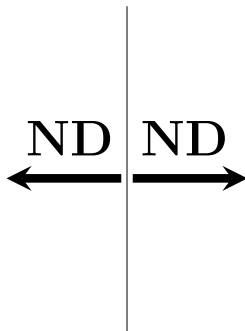
Lemma (Cyr & Kra 2015)

$c$  of low complexity and uniformly recurrent.

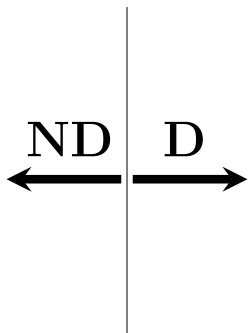
$\exists \mathbf{u}, \overline{\mathcal{O}(c)}$  non-deterministic for  $\mathbf{u}$  and  $-\mathbf{u}$ ,

$\Rightarrow$

$c$  is periodic



# One-sided Determinism



Last case: **One-sided Determinism**

.. can be eliminated

Theorem (Kari & M. 2020)

$c$  2D with non-trivial annihilator.

$\exists d \in \overline{\mathcal{O}(c)}$  such that  $\overline{\mathcal{O}(d)}$  **has no direction of one-sided determinism.**

# Main Theorem

Theorem (Kari & M. 2020)

$c$  such that  $\exists m, n, P_c(m, n) \leq mn$ ,

$\exists d \in \overline{\mathcal{O}(c)}$  which is periodic

Proof sketch:

- “Eliminate” all one-sided deterministic directions  $\rightarrow d \in \overline{\mathcal{O}(c)}$
- $d$  is periodic ! ([Boyle & Lind 1997] + [Cyr & Kra, 2015])

# Consequence 1: Uniform Recurrence

Corollary (Kari & M. 2020)

Nivat's conjecture holds for  $c$  uniformly recurrent

Proof.

$$c \in \mathcal{A}^{\mathbb{Z}^2}, P_c(m, n) \leq mn.$$

There exists  $d \in \overline{\mathcal{O}(c)}$  periodic.

$$\overline{\mathcal{O}(c)} \text{ uniformly recurrent} \Rightarrow \overline{\mathcal{O}(c)} = \overline{\mathcal{O}(d)}.$$

$\overline{\mathcal{O}(d)}$  contains only periodic configurations  $\Rightarrow c$  periodic. □

What's next ?

# Nivat's Conjecture

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Non uniformly recurrent configurations: contains arbitrarily large periodic regions

What is the geometry of these regions?

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Non uniformly recurrent configurations: contains arbitrarily large periodic regions

What is the geometry of these regions?

→ What complexity can have aperiodic SFTs?



# What complexity can have aperiodic SFTs?

Theorem (Kari & M. 2020)

For all **aperiodic** subshifts  $X$ ,  $\forall c \in X, \forall m, n \in \mathbb{N}$ ,

$$P_c(m, n) > mn$$

# What complexity can have aperiodic SFTs?

$\forall X, \forall c \in X, c$  periodic

$\forall X, \exists c \in X, c$  periodic

$\exists X, \forall c \in X, c$  not periodic



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What is the smallest possible  $C$ ? (we know  $C > 0$ )

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Conjecture

$C > 1$

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Theorem (M. & Petit-Jean 2021)

$\sigma : \mathcal{A} \rightarrow \mathcal{A}^{\{1, \dots, n\}^2}$  primitive substitution with **determining position**. If  $X^\sigma$  is aperiodic, then  $\exists C > 1, \forall c \in X^\sigma, \forall n,$

$$P_c(n, n) \geq Cn^2$$

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Conjecture

True without the *determining position* assumption

Thank you !