p-adic construction of CM-curves

David Kohel & Christophe Ritzenthaler

Abstract

We give a *p*-adic analytic construction of the invariants of CM curves of genus 2, obtained by the 2-adic AGM lifting algorithm. This construction provides an alternative to the complex analytic approach for reconstructing the invariants of curves. By reduction modulo a suitable large prime, the CM invariants of these curves enable the efficient construction of curves of with known group order suitable for cryptosystems based on the discrete logarithm problem.

1 Introduction

The traditional approach to CM constructions in genus 1 has been through evaluate of the j-function on an upper half complex plane at special points corresponding to lattices with complex multiplication [4] or using special modular functions of higher level as in Yui and Zagier [17] or Enge and Morain [2]. This construction has been extends to genus 2 curves, using theta functions on Siegel upper half plane (see, e.g., van Wamelen [14] and Weng [16]).

The *p*-adic point counting algorithms of Satoh and generalizations such as Mestre's AGM method determine the number of points on an elliptic or hyperelliptic curves by constructing a *p*-adic canonical lift. Although conceived for the purpose of point counting these algorithms are in fact *p*-adic analytic analogues of the complex analytic CM constructions cited above. Couveignes and Henocq [1] developed the theory of this method when applied to the *j*-function in genus 1.

In the present work, we utilise the AGM construction for genus 2 curves to lift invariants of a hyperelliptic curve over a finite field of characteristic 2 to an extension of \mathbb{Q}_2 , then use lattice reduction to reconstruct the minimal polynomials of these invariants over \mathbb{Q} . The algorithm uses only the elementary recursive construction of the AGM algorithm, applied to curves over small finite fields, together with LLL reduction to rationally reconstruct the invariants.

2 Canonical Lift by the AGM

We recall in this section the principle and the formulas of the AGM algorithm for genus 2 curves. For proofs we refer to Lercier and Lubicz [5], Mestre [6] or Ritzenthaler [12]. Let $q = 2^n$, set $k = \mathbb{F}_q$, let $K = \mathbb{Q}_q$ be the unramified extension of \mathbb{Q}_2 of degree n, and let \mathbb{Z}_q be its ring of integers. Then the Galois group $\operatorname{Gal}(K/\mathbb{Q}_2)$ is generated by the Frobenius automorphism which we denote by σ .

The AGM algorithm applies to any ordinary hyperelliptic curve \tilde{C} , which we may represent in Weierstrass form:

$$C/k: y^2 + \tilde{v}(x)y = \tilde{u}(x)\tilde{v}(x)$$
(1)

where \tilde{v} and \tilde{u} are degree 3 monic polynomials such that \tilde{v} is square-free. Note that for any curve

$$\tilde{C}/k: y^2 + \tilde{v}(x)y = f(x)$$

such that (v, f) = w, we can set $c = 1/\sqrt{f/w} \mod (v/w)$, and set make a change of variables $y \mapsto y + c$ to put \tilde{C} in the form (1).

Such a curve \tilde{C} is a genus 2 curve and is ordinary, i.e the Jacobian \tilde{J} of \tilde{C} , has four 2torsion points defined over some extension field. We know then that there exists a principally polarized abelian surface $(J, \lambda)/K$ which lifts the principally polarized Jacobian $(\tilde{J}, \tilde{\lambda})/k$ together with its ring of endomorphisms: $\operatorname{End}_K(J) \simeq \operatorname{End}_k(\tilde{J})$.

Using the AGM algorithm, we can construct sequences of 2-adic numbers which converge 2-adicly to 'invariants' associated to (J, λ) . This is achieved by the following process :

- 1. Replace k by a finite extension (of degree up to three) such that the roots of \tilde{v} are defined.
- 2. Lift \tilde{C} over K: Lift \tilde{v} and \tilde{u} to v(x) and u(x) in K[x] and then let

$$C/K: Y^2 = (2y + v(x))^2 = v(x)(v(x) + 4u(x)).$$

Since \tilde{v} splits in k with distinct roots, we can write in K,

$$C/K: Y^2 = \prod_{i=1}^3 (x - x_i) \prod_{i=1}^3 (x - (x_i + 4s_i)).$$

3. Initialization of theta characteristics: Denote by

$$e_1 = x_1,$$
 $e_3 = x_2,$ $e_5 = x_3,$
 $e_2 = x_1 + 4s_1,$ $e_4 = x_2 + 4s_2,$ $e_6 = x_3 + 4s_3$

The Thomae formulas give us 4 initial invariants

$$A = (e_1 - e_3)(e_3 - e_5)(e_5 - e_1)(e_2 - e_4)(e_4 - e_6)(e_6 - e_2)$$

$$B = (e_1 - e_3)(e_3 - e_6)(e_6 - e_1)(e_2 - e_4)(e_4 - e_5)(e_5 - e_2)$$

$$C = (e_1 - e_4)(e_4 - e_5)(e_5 - e_1)(e_2 - e_3)(e_3 - e_6)(e_6 - e_2)$$

$$D = (e_1 - e_4)(e_4 - e_6)(e_6 - e_1)(e_2 - e_3)(e_3 - e_5)(e_5 - e_2)$$

We recall that these numbers are 2-adic analogs of the respective complex values :

$$\vartheta^{[00]}_{00}(0)^4, \ \vartheta^{[00]}_{10}(0)^4, \ \vartheta^{[00]}_{01}(0)^4, \ \vartheta^{[00]}_{11}(0)^4$$

We initialize $(A_0, B_0, C_0, D_0) := (1, \sqrt{B/A}, \sqrt{C/A}, \sqrt{D/A})$, where the square root of an element of the form $1 + 8\mathbb{Z}_q$ is taken as the unique element of \mathbb{Z}_q of the form $1 + 4\mathbb{Z}_q$.

4. Lifting process: We use the duplication formula to obtain a 4-tuple of invariants

$$(A_n, B_n, C_n, D_n)$$

as elements of \mathbb{Z}_q :

$$A_{n+1} = \frac{A_n + B_n + C_n + D_n}{\sqrt{A_n B_n} + \sqrt{C_n D_n}} \quad C_{n+1} = \frac{\sqrt{A_n C_n} + \sqrt{B_n D_n}}{\sqrt{A_n D_n} + \sqrt{B_n C_n}}$$
$$B_{n+1} = \frac{\sqrt{A_n B_n} + \sqrt{C_n D_n}}{2} \quad D_{n+1} = \frac{\sqrt{A_n D_n} + \sqrt{B_n C_n}}{2}$$

These invariants do not converge but if we denote the invariants associated to J by $(A_{\infty}, B_{\infty}, C_{\infty}, D_{\infty})$, we have

$$(A_n, B_n, C_n, D_n) \equiv (A_\infty, B_\infty, C_\infty, D_\infty)^{\sigma^n} \mod 2^n,$$

where σ is the Frobenius automorphism of $\mathbb{Z}_q/$. In particular, each of the sequences of invariants

$$(A_{kr+i}, B_{kr+i}, C_{kr+i}, D_{kr+i}),$$

for fixed i in $1 \le i \le r$, does converge as k goes to infinity. But since we may consider any of the Galois conjugates of Igusa invariants, we terminate the algorithm at any step n to obtain a precision of n bits.

Finally, we note that the algorithmic improvements of Lercier and Lubicz [5] to obtain quadratic convergence is applicable here.

3 Computation of the *p*-adic invariants

The sequence of values A_n , B_n , C_n , D_n of the preceding section describe a cycle of Galois conjugate invariants of the canonical lift (J, λ) to K of our original Jacobian $(\tilde{J}, \tilde{\lambda})$ over k. In genus 2 the canonical lift is itself the Jacobian of a genus 2 curve C over K. We now describe how to determine the invariants of the curve C/K, from a set of invariants A_n , B_n , C_n , and D_n .

We proceed in two steps as described by van Wamelen [14]. Recall that over \mathbb{C} , if C is given by the Rosenhain normal form

$$C: y^2 = x(x-1)(x-\lambda_1)(x-\lambda_2)(x-\lambda_3)$$

then the λ_i are given by the following expressions :

$$\lambda_1 = -\frac{\vartheta_1^2 \vartheta_3^2}{\vartheta_6^2 \vartheta_4^2}, \quad \lambda_2 = -\frac{\vartheta_2^2 \vartheta_3^2}{\vartheta_6^2 \vartheta_5^2}, \quad \lambda_3 = -\frac{\vartheta_2^2 \vartheta_1^2}{\vartheta_4^2 \vartheta_5^2}$$

where

$$\vartheta_1 = \vartheta_{10}^{[00]}(0), \quad \vartheta_2 = \vartheta_{11}^{[00]}(0), \quad \vartheta_3 = \vartheta_{10}^{[01]}(0),$$

$$\vartheta_4 = \vartheta_{00}^{[10]}(0), \quad \vartheta_5 = \vartheta_{01}^{[10]}(0), \quad \vartheta_6 = \vartheta_{00}^{[11]}(0).$$

We then use 2-adic analogues that we can compute by means of the general duplication formulas (see Mumford [10], and [11]), namely we set

$$\begin{split} \vartheta_1^2 &= B_n, & \vartheta_2^2 &= D_n, \\ \vartheta_3^2 &= \frac{\sqrt{A_{n-1}B_{n-1}} - \sqrt{C_{n-1}D_{n-1}}}{2}, & \vartheta_4^2 &= \frac{A_{n-1} - B_{n-1} + C_{n-1} - D_{n-1}}{4}, \\ \vartheta_5^2 &= \frac{\sqrt{A_{n-1}C_{n-1}} - \sqrt{B_{n-1}D_{n-1}}}{2}, & \vartheta_6^2 &= \frac{A_{n-1} - B_{n-1} - C_{n-1} + D_{n-1}}{2}. \end{split}$$

Given λ_i we can then compute the Igusa invariants I_2 , I_4 , I_6 , and I_{10} (for details we refer to van Wamelen [14]) and define the absolute invariants

$$i_1 = I_2^5/I_{10}, \ i_2 = I_2^3I_4/I_{10}, \ i_3 = I_2^2I_6/I_{10}.$$

4 Rational reconstruction of the invariants

From the *p*-adic invariants it remains to determine a set of defining relations over \mathbb{Z} . For this purpose it is desirable to predetermine the degree of relations among the absolute invariants. This degree can be explicitly determined, from the data of a CM type for K. However, in the case that K/Q defines a cyclic or non-normal quartic extension, that the totally real subfield L of K has class number one, and the only roots of unity in K are $\{\pm 1\}$, we have the following theorem of Weng [15, Theorem 3.1].

Theorem 1. The number of classes of Igusa invariants of a CM type for the maximal order of K equals the class number h_K of K if K is cyclic and $2h_K$ if K is a non-normal quartic extension.

From these absolute invariants, we use LLL on the space of *p*-adic relations among the powers $1, i_k, i_k^2 \dots, i_k^n$ of degree *n* to solve for

$$H_1(i_1) = H_2(i_2) = H_3(i_3) = 0.$$
 (2)

Such relations appear as short vectors in the space of all relations over \mathbb{Z}_p to some precision p^N . In addition, we reconstruct additional relations

$$L_1(i_1, i_2, i_3) = L_2(i_1, i_2, i_3) = 0, (3)$$

in order to record the dependencies among the different invariants. This removes the problem of combinatorial matching of up to n^3 possible combinations of roots over some finite field \mathbb{F}_p .

We note that the polynomials H_1 , H_2 , and H_3 are not in general monic. The possible prime divisors of the leading coefficient are characterised by Goren and Lauter [9]. Although the exact powers of these leading coefficients are not known, it is possible to clear denominators in the absolute invariants and reconstruct first the leading coefficients using a much smaller precision.

A more critical issue is the identification of a representative curve whose Jacobian has maximal endomorphism ring. It is necessary to have a mechanism to distinguish and discard curves associated to the nonmaximal orders. The following theorem provides such as test.

Theorem 2. Let χ be the minimal polynomial of the Frobenius endomorphism Frob_q on the Jacobian J of a genus 2 curve C/\mathbb{F}_q . Let π be any root of this polynomial and set $K = \mathbb{Q}(\pi)$ and $\overline{\pi} = q/\pi$. If the set

$$\left\{\frac{f_1(\pi)}{m_1},\ldots,\frac{f_t(\pi)}{m_t}\right\}$$

for $(m_i, q) = 1$ generates the maximal order O_K over $\mathbb{Z}[\pi, \overline{\pi}]$, then $\operatorname{End}(J) = O_K$ if and only if $f_i(\operatorname{Frob}_q)$ is the zero map on $J[m_i]$ for all *i*.

N.B In practice it suffices to check only for each maximal prime power $p_i^{e_i}$ dividing each m_i .

5 Algorithm and Examples

Strategy:

1. For a given field $k = \mathbb{F}_{2^n}$, choose curve defined by u, v in k[x], hence with field of moduli equal to k, then determine theta constants over some extension.

- 2. Determine the index of $\mathbb{Z}[\pi, \bar{\pi}]$ in the maximal order O_K , and the group structure of quotient $A = O_K / \mathbb{Z}[\pi, \bar{\pi}]$.
- 3. Let $f_1(\pi)/m_1, \ldots, f_t(\pi)/m_t$ generate O_K over $\mathbb{Z}[\pi, \overline{\pi}]$. For each m_i determine the action of π on $J[m_i]$ and reject the curve if the restriction of $f_i(\pi)$ to $J[m_i]$ is nonzero.
- 4. Lift the theta constants and reconstruct by LLL the defining relations for the CM igusa invariants defining those curves whose Jacobian J has O_K embedded in End(J).

Note that by choosing a curve over its field of moduli, rather than the extension field over which the Weierstrass points are defined, we select a curve whose Jacobian is more likely to be in the class of the maximal endomorphism ring. Such a curve minimizes both the degree and the size of coefficients in the relations for the Igusa invariants.

Examples. Here we provide a few examples of canonical lifts of the Igusa invariants of hyperelliptic curves of the form

$$C: y^2 + v(x)y = v(x)u(x)/\mathbb{F}_{2^n},$$

and their application to explicit constructions of Jacobians suitable for cryptography.

1. For the curve C/\mathbb{F}_2 with $v = x^3 + 1$ and $u = x^2$, we find relations for the canonical lifts fo the Igusa invariants:

$$\begin{split} &i_1^2 - 531441i_1 + 55788550416, \\ &i_2^2 - 426465i_2 - 68874753600, \\ &i_3^2 - 216513i_3 - 221011431552, \\ &140i_1 - 243i_2 + 135i_3, \\ &69i_1 - 119i_2 + 66i_3 - 104976. \end{split}$$

The minimal polynomial of Frobenius in End(J) is equal to

$$x^4 + 2x^3 + 3x^2 + 4x + 4,$$

defining an imaginary quadratic extension of the real quadratic field $\mathbb{Q}(\sqrt{2})$.

2. For the curve C/\mathbb{F}_2 with $v = x^3 + x^2 + 1$ and $u = x^2 + 1$ we find relations for the canonical lifts for the Igusa invariants:

$$\begin{array}{l} 4i_1^2 + 8218017i_1 + 146211169851,\\ i_2^2 + 1008855i_2 - 342014432400,\\ i_3^2 + 1368387i_3 - 240090131376,\\ 4480i_1 + 7499i_2 - 12255i_3,\\ 716i_1 + 1212i_2 - 1971i_3 - 1666737\end{array}$$

The minimal polynomial of Frobenius in End(J) is equal to

$$x^4 + x^3 + x^2 + 2x + 4$$

defining an imaginary quadratic extension of the real quadratic field $\mathbb{Q}(\sqrt{13})$.

3. For the curve C/\mathbb{F}_2 with $v = x^3 + x^2 + 1$ and $u = x^2$ we find relations for the canonical lifts for the Igusa invariants:

 $\begin{array}{l} 4i_1^2+115322697i_1-10896201253125,\\ i_2^2+9073863i_2-2152336050000,\\ i_3^2+14410143i_3-1214874126000,\\ 896i_1+369i_2-2025i_3,\\ 300i_1+122i_2-677i_3+273375 \end{array}$

The minimal polynomial of Frobenius in End(J) is equal to

$$x^4 + x^3 + 3x^2 + 2x + 4,$$

defining an imaginary quadratic extension of the real quadratic field $\mathbb{Q}(\sqrt{5})$.

6 Conclusion

The AGM provides a relatively elementary and effective alternative to the complex analytic construction of complex multiplication for genus 2 curves. We note that not all CM orders arise in this way, but those curves that do have good reduction at 2 and small class number appear among the curves over small fields \mathbb{F}_2^n . The approach through *p*-adic lifting also permits us to treat curves whose Jacobians are not absolutely simple. In order to capture additional orders, corresponding to Jacobians with bad reduction at 2, it would be desirable to extend the algorithmic theory of canonical lifts to curves of genus 2 in odd characteristic.

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Example Finale. Let $C: y^2 + v(x)y = u(x)v(x)$ be the hyperelliptic curve over $\mathbb{F}_{2^3} = \mathbb{F}_2[w]$ where $w^3 + w + 1 = 0$ where u and v are given by

$$u = (w^{2} + w + 1)x^{2} + w^{2}x + w^{2},$$

$$v = x^{3} + (w^{2} + w + 1)x^{2} + x + w + 1.$$

The minimal polynomial of Frobenius on the Jacobian of C is

 $x^4 - 3x^3 + 3x^2 - 24x + 64,$

defining an imaginary quadratic extension of the real quadratic field $\mathbb{Q}(\sqrt{61})$. The defining relations of canonical lifts of the Igusa invariants are given below.

 $2^{6}3^{42}i_{1}^{6} - 2344912105503116116288576047953057125392i_{1}^{5}$ $-112639584390304238456172276845130150039402556586283156i_1^4$ $-2177415103395854060041246748534717663224784831560700934285483051075i_1^3$ $-1593641994054440870937630653070363836936366222692321471303808012543988702i_1^2$ + 32299720850335379144290409627740329840675572467939277123595091705537581712591977043, $3^{18}i_2^6 + 30345890982308051019805350i_2^5$ $-288136191649832893917062077388710908375i_2^4$ $+753110832515821367749096990899427029369367852656375i_{2}^{3}$ $-649127309475920539312400482687597914255658885551562830000i_{2}^{2}$ $+ 512065244591992233358858681228726038539915018527646447680800000 i_2$ -242729201551569096286616270971131120449527443900342023922233408000000, $3^{24}i_3^6 + 27437461181384763694011881346i_3^5$ $-352040806049318452655962733807057489240331i_3^4$ $+ 1178922153334081066484173968480725700444739639422966003i_3^3$ $+509928790982645514856427558535377505816658890920020722687216i_3^2$ $+22813028282617457487855156583191936594982551082177632973015943424i_3$ -194627707132727224036285973133204401034007902817343828521298858611945472, $633895738920000i_1^3 + 8517595035131037i_1^2i_2 - 2422318926838275i_1^2i_3$ $+528887012556497760i_1^2-2671415018933342i_1i_2^2+10103099744994882i_1i_2i_3$ $+ 11002415784338674i_2^3 - 16195247750833904i_2^2i_3 + 800164846490774071i_2^2i_3 + 800164846490i_3 + 8001648i_3 + 8001648i_3 + 8001648i_3 + 8001648i_3 + 8001648i_3 + 8000i_3 + 800i_3 +$ $+228622640238253145i_2i_3,$ $52586040050922240i_1^3 + 348046133200631478i_1^2i_2 + 19788972081057810i_1^2i_3$ $-1585090558318459827i_3^3 - 10377834109186130040i_3^2 - 12385238120639343570i_3$ $+ 3160028075123540i_2^3 - 19415412647408141i_2^2i_3 - 11227855503503951i_2^2$ $+42818455041104040i_{3}^{2}$