Construction of CM moduli by *p*-adic lifting

David R. Kohel The University of Sydney

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Canonical lifts of elliptic curves Canonical lifting algorithm Canonical lifting algorithm Canonical lifts of abelian surfaces Constructive CM algorithms for genus 2 Effective Class Field Theory Cryptographic applications Database of CM moduli

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(i) $\tilde{A}/R \longleftarrow A/\mathbb{F}_q$ and (ii) $\operatorname{End}_R(\tilde{A}) = \operatorname{End}_{\mathbb{F}_q}(A)$.

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In order to understand canonical lifts, we need also to understand the endomorphism rings of such curves. In this one-dimensional case, the only CM endomorphism rings are orders in an imaginary quadratic field K.

Example of a canonical lift

Example. The simplest example of such a curve is

$$E/\mathbb{F}_p: y^2 = x^3 - x,$$

where $p = 1 \mod 4$, which has canonical lift \tilde{E}/\mathbb{Z}_p : $y^2 = x^3 - x$.

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The objective of our investigation, however, is to recover the *invariants* or *moduli* of \tilde{E} , which in the elliptic curve case is the *j*-invariant. For p = 17 we compute $j(E) = 18 \in \mathbb{F}_{19}$, however the canonical lift has invariant $j(\tilde{E}) = 12^3 \in \mathbb{Z}$.

Example of a canonical lift

N.B. The *j*-invariant of the canonical lift of E/\mathbb{F}_p lies in \mathbb{Z}_p , but is algebraic over \mathbb{Z} ,

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N.B. The *j*-invariant of the canonical lift of E/\mathbb{F}_p lies in \mathbb{Z}_p , but is algebraic over \mathbb{Z} , and moreover generates the Hilbert class field over $\mathcal{K} = \operatorname{End}(E) \otimes \mathbb{Q}$.

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By lifting to sufficient precision we verify that $j = j(\tilde{E})$ satisfies the quadratic relation:

$$j^2 + 191025j - 121287375.$$

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A *p*-adic algorithm for constructive CM must

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The first step replaces the p-adic numbers with complex numbers in analogous analytic constructions. Rather than a period lattice, the input is a suitable curve which we lift p-adically.

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A prior algorithm for constructing canonical lifts for point counting was developed by Satoh. Efficient versions were introduced by Mestre, which generalise to higher dimension.

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Canonical lifts of abelian surfaces

In higher dimension, we first need explicit models for abelian varieties, secondly, and explicit descriptions of their invariants or moduli. For genus 2 curves "most" abelian surfaces are Jacobians of genus 2 curves, and we have an explicit algebraic description of their invariants by Igusa (following analytic invariants of Clebsch in the 19th century). In the above construction we replace j with a triple of Igusa invariants (j_1, j_2, j_3) on \mathcal{M}_2 , and find suitable correspondences relating the invariants of (p, p)-isogenous abelian varieties.

Constructive CM algorithms for genus 2

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- ▶ *p*-adic lifting of (ℓ, ℓ) -isogenies (K., adapting above to $p \neq \ell$). The first uses Richelot isogenies between Jacobians of curves in Rosenhain form:

$$y^2 = x(x-1)(x-\lambda_1)(x-\lambda_2)(x-\lambda_3),$$

and the 3-adic algorithm makes use of correspondence equations of algebraic theta functions.

Example of genus 2 CM construction

Example. Let C be defined over \mathbb{F}_2 with model

$$y^{2} + (x^{3} + 1)y = x(x^{3} + 1).$$

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Its Jacobian is an abelian surface with complex multiplication by the maximal order of the number field

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Then canonically lifted Igusa invariants (j_1, j_2, j_3) satisfy:

$$\begin{array}{ll} j_1^2-531441 j_1+55788550416, & 34 j_2-36864 j_3+10206, \\ 8 j_2^2-4374 j_2-9565938, & j_1+176 j_2-73728 j_3-27483. \\ 8 192 j_3^2-8667 j_3-6561, & \end{array}$$

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We want to make use of this knowledge...

Effective Class Field Theory

Returning to the previous example, we find that the reflex field is

$$\mathcal{K}^r = \mathbb{Q}[x]/(x^4 + 5x^2 + 2) \cong \mathbb{Q}\left(i\sqrt{(5+\sqrt{17})/2}\right)$$

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With Claus Fieker, we are combining algorithms for effective class field theory, to determine H, with the algebraic reconstruction of (j_1, j_2, j_3) , to determine them as elements of the known field H.

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Cryptographic applications

Example. Let C be the curve $y^2 + h(x)y = f(x)$ over

$$\mathbb{F}_8 = \mathbb{F}_2[t]/(t^3+t+1),$$

with h(x) = x(x + 1) and $f(x) = x(x + 1)(x^3 + x^2 + t^2x + t^3)$.

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The field K has class number is 3, and there exist 6 isomorphism classes of principally polarized abelian varieties.

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We can construct the defining ideal of relations in Igusa invariants (j_1, j_2, j_3) from the canonical lift of (the Jacobian of) C.

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For example, the invariant j_1 satisfies a minimal polynomial:

 $H_1(X) = 2^{18} 5^{36} 7^{24} X^6$

 $-\,11187730399273689774009740470140169672902905436515808105468750000\,X^{5}$

 $+ \ 501512527690591679504420832767471421512684501403834547644662988263671875000 \ X^4$

 $- \ 1011 \ 24 \ 092 \ 427 \ 87 \ 391 \ 78 \ 66 \ 76 \ 28 \ 46 \ 33 \ 73 \ 05 \ 75 \ 04 \ 76 \ 14 \ 54 \ 31 \ 35 \ 57 \ 20 \ 25 \ 66 \ 74 \ 68 \ 22 \ 14 \ 32 \ 70 \ 42 \ 6 \ 38 \ 57 \ 80 \ 82 \ 62 \ 92 \ 3 \ X^3$

+ 118287000250588667564540744739406154398135978447792771928535541240797386992091828213521875 χ^{2}

 $-\,{}^{2^1}{}^{5^0}{}^{5^{10}}{}^{11^1}{}^{1}{}^{3^1}{}^{53^1}{}^{701^1}{}^{1}{}^{6319^1}{}^{69938793494948953569198870004032131926868578084899317}\, X$

 $+\ 3^{60} 5^{15} 23^5 409^5 179364113^5$

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Choosing the 120-bit prime

p = 954090659715830612807582649452910809,

and solving a norm equation in the endomorphism ring $\mathcal{O}_{\mathcal{K}}$,

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and solving a norm equation in the endomorphism ring \mathcal{O}_K , we determine that the Jacobian of some curve over \mathbb{F}_p with CM by \mathcal{O}_K will have prime order

 $\begin{array}{l}910288986956988885753118558284481029 \\311411128276048027584310525408884449.\end{array}$

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$$\begin{split} C: y^2 &= x^6 + 827864728926129278937584622188769650 \, x^4 \\ &\quad + 102877610579816483342116736180407060 \, x^3 \\ &\quad + 335099510136640078379392471445640199 \, x^2 \\ &\quad + 351831044709132324687022261714141411 \, x \\ &\quad + 274535330436225557527308493450553085. \end{split}$$

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A test of a random point on the Jacobian verifies the group order.

Database of CM moduli

A comprehensive database for CM invariants in genera 1 and 2 is being developed:

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providing an interface for the interrelated invariants of CM fields K, their Hilbert class fields, and CM moduli of abelian varieties.