## Mahler equations for Zeckendorf numeration

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Combinatorics on Words Seminar October 2024 Prologue: numeration systems and automaticity

Act 1: the classical base-q numeration

Act 2: the Zeckendorf numeration



## Outline

### Prologue: numeration systems and automaticity

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Act 1: the classical base-q numeration

Act 2: the Zeckendorf numeration

## Numeration systems

A numeration system (U, B) consists of

- ▶ a sequence of natural numbers  $U = (u_n)_{n \ge 0}$  with  $u_0 = 1$ ,
- ▶ a finite ordered digit set B such that for each  $n \in \mathbb{N}$  there are  $b_k, \ldots, b_0 \in B$  with  $n = \sum_{i=0}^k b_i u_i$ .

We say that  $b_k \cdots b_0$  is the canonical representation of n if  $b_k \cdots b_0$  is the greatest representation of n for the lexicographic order.

We write

$$(n)_U := b_k \cdots b_0$$
 and  $[b_k \cdots b_0]_U := n$ 

Example (The base-q numeration system for  $q \ge 2$ )  $U = (q^n)_{n \ge 0}, B = \{0, 1, \dots, q-1\}$ . If q = 3, then  $(26)_3 = 222$ .

# Zeckendorf numeration

Recall the Fibonacci numbers  $(F_n)_{n \ge 0}$ , defined by

$$\begin{split} F_{-2} &= 0, \\ F_{-1} &= 1, \\ F_n &= F_{n-1} + F_{n-2} \text{ for } n \geqslant 0. \end{split}$$

The Zeckendorf numeration system is  $Z = ((F_n)_n, B = \{0, 1\}).$ 

The canonical expansion  $(n)_Z = b_n \cdots b_0$  satisfies  $b_i b_{i+1} = 0$  for each *i*.

### Example

	n	55	34	21	13	8	5	3	2	1	
-	61	0	1	0	1	1	1	0	0	1	
	61	0	1	1	0	0	1	0	0	1	
	61	1	0	0	0	0	1	0	0	1	
(6)	$1)_Z$	= 10	0000	1001	and	[10	111	100	$l]_Z$	= 6	;]

# Automaticity in U

A sequence  $(a_n)_{n\geq 0}$  taking values in a finite set A is *U*-automatic if there is a deterministic finite automaton whose output is  $a_n$  when fed with  $(n)_U$ . If U is the base-q numeration, we will say that  $(a_n)_{n\geq 0}$  is q-automatic.

Example (Thue-Morse  $a_n = |(n)_2|_1 \mod 2$ )



 $(17)_2 = 10001$  and  $s \xrightarrow{1}{\rightarrow} t \xrightarrow{0}{\rightarrow} t \xrightarrow{0}{\rightarrow} t \xrightarrow{1}{\rightarrow} s$  and  $s \xrightarrow{0}{\rightarrow}$ 

Thus  $a_{17} = 0$ .  $a_n = 0$  precisely when  $(n)_2$  contains an even number of the digit 1.

### Prologue: numeration systems and automaticity

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Characterisation of q-automaticity for  $q = p^n$ , p prime

Theorem (Christol 1980)

Let  $(a_n)_{n\geq 0}$  be a sequence in  $\mathbb{F}_q$ , with  $q = p^n$  for some  $n \geq 1$ . Then  $(a_n)_{n\geq 0}$  is q-automatic if and only if  $f(x) = \sum_{n\geq 0} a_n x^n$  is algebraic over  $\mathbb{F}_q(x)$ .

Example (Catalan numbers  $C_n = \frac{1}{n+1} \binom{2n}{n}$ )

• 
$$y = \sum_{n \ge 0} C_n x^n$$
 satisfies  $xy^2 - y + 1 = 0$  over  $\mathbb{Q}$ ,

•  $y = \sum_{n \ge 0} (C_n \mod 3) x^n$  satisfies  $xy^2 + 2y + 1 = 0$  over  $\mathbb{F}_3$ , and hence  $(C_n \mod 3)_{n \ge 0}$  is automatic.



# Limitations of Christol's theorem

Christol's theorem does not provide a characterisation of q-automatic sequences if q not a power of a prime.

Question Is there a generalisation of Christol's theorem to all *q*-automatic sequences?

Answer Yes.

Question Is there a version of Christol's theorem for the Zeckendorf numeration system?

Answer Yes!

## Christol's theorem, example

Let us find the annihilating polynomial for the Thue-Morse sequence, defined as

 $a_n = 0$  precisely when  $(n)_2$  contains an even number 1s.

$$f(x) = \sum_{n} a_{2n} x^{2n} + \sum_{n} a_{2n+1} x^{2n+1}$$
$$= \sum_{n} a_n x^{2n} + x \sum_{n} (a_n + 1) x^{2n}$$

so, with  $s(x) = \frac{1}{1+x}$ ,  $f(x) = (1+x)f(x^2) + xs(x^2) \& s(x) = (1+x)s(x^2)$  (1) and

$$f(x^2) = (1+x^2)f(x^4) + x^2s(x^4), & s(x^2) = (1+x^2)s(x^4)$$

Substituting (2) in (1) we get that f(x) is a root of the Ore polynomial

$$xf(x) = (1+x)f(x^2) + (1+x)^4 f(x^4).$$

From automatic to regular sequences

### Definition

A sequence  $(a_n)_{n \ge 0}$  taking values in a finite set A is *U*-automatic if there is a deterministic finite automaton whose output is  $a_n$  when fed with  $(n)_U$ .

### Definition (Allouche-Shallit)

A sequence  $(a_n)_{n\geq 0}$  taking values in a finite set A commutative ring R is U-automatic U-regular if there is a deterministic finite automaton weighted automaton whose output is  $a_n$  when fed with  $(n)_U$ .

## Theorem (Allouche-Shallit 1992)

A sequence is q-regular and takes on finitely many values if and only if it is q-automatic.

# Examples

All from Allouche-Shallit's article, 1992:

- ▶  $a_n = #$  1's in  $(n)_2$  defines a 2-regular sequence
- ▶ the sequence

$$0, 2, 6, 8, 20, 24, \ldots,$$

which lists the numerators of the left endpoints of the Cantor set, is 2-regular.

► 
$$a_n = (n^j)_{n \ge 0}$$
 is 2-regular,

► 
$$a_n = \sum_{i=1}^n \lfloor \log_a i \rfloor$$
 is *q*-regular.

- The number of comparisons required to mergesort n items,
- For  $a \in \mathbb{R}$ ,  $(a^n)_{n \ge 0}$  is q-regular if and only if a = 0 or a is a root of unity.

# Weighted automata

A weighted automaton  $\mathcal{A}$  with weights in a ring R consists of

- $\blacktriangleright$  a finite state set S,
- ▶ an alphabet B
- ▶ a transition weight function  $\Delta : S \times B \times S \to R$  which assigns a weight to each labelled edge, denoted  $s \xrightarrow{b:r} s'$ , and
- ▶ initial and final weight functions  $I: S \to \mathbb{R}$  and  $F: S \to \mathbb{R}$ .

### Example

Let  $B = \{0, 1\}$  and  $R = \mathbb{F}_2$ .



### Generating sequences using weighted automata

In a weighted automaton, there may be many paths that a given word labels. We are interested in the sum of the weights of all paths that this word labels.

The word 10110 labels three different paths, each of weight 1:

$$\begin{array}{c} \xrightarrow{1} s \xrightarrow{1:1} t \xrightarrow{0:1} t \xrightarrow{1:1} t \xrightarrow{1:1} t \xrightarrow{0:1} t \xrightarrow{1} \\ \xrightarrow{1} s \xrightarrow{1:1} s \xrightarrow{0:1} s \xrightarrow{1:1} t \xrightarrow{1:1} t \xrightarrow{0:1} t \xrightarrow{1} \\ \xrightarrow{1} s \xrightarrow{1:1} s \xrightarrow{0:1} s \xrightarrow{1:1} s \xrightarrow{1:1} t \xrightarrow{0:1} t \xrightarrow{1} \end{array}$$

and since  $(22)_2 = 10110$  and  $R = \mathbb{F}_2$ , then  $u_{22} = 3 \mod 2 = 1$ .



## q-Mahler equations

Let R be any commutative ring and let  $q \ge 2$ . Define the linear operator  $\Phi: R[\![x]\!] \to R[\![x]\!]$  as

$$\Phi(f(x)) = f(x^q).$$

Let  $A_i(x) \in R[x]$  be polynomials. The equation

$$P(x,y) = \sum_{i=0}^{d} A_i(x)\Phi^i(y) = 0$$

is called a *q*-Mahler equation. If  $f \in R[\![x]\!]$  satisfies P(x, f(x)) = 0, then it is called *q*-Mahler.

If  $q = p^k$ , then a q-Mahler equation over  $\mathbb{F}_q$  is just a polynomial:

$$\left(\sum_{n\geq 0} f_n x^n\right)^q = \sum_{n\geq 0} f_n^q x^{qn} = \sum_{n\geq 0} f_n x^{qn}$$

#### Theorem (Christol 1980)

Let q be a power of a prime, and let  $(u_n)_{n \ge 0}$  be a sequence over  $\mathbb{F}_q$ . The  $(u_n)_{n \ge 0}$  is q-automatic if and only if  $f(x) = \sum_{n \ge 0} a_n x^n$  is algebraic over  $\mathbb{F}_q(x)$ .

#### Theorem (Becker 1992, Dumas 1993)

Let  $q \ge 2$ , and let  $(u_n)_{n \ge 0}$  be a sequence over a commutative ring R.

- ▶ If  $(u_n)_{n \ge 0}$  is q-regular sequence then  $f(x) = \sum_{n \ge 0} a_n x^n$  is the solution of a q-Mahler equation,
- if  $f(x) = \sum_{n \ge 0} a_n x^n$  is the solution of an isolating q-Mahler equation, i.e., of the form  $y = \sum_{i=1}^d A_i(x) \Phi^i(y)$ , then  $(u_n)_{n \ge 0}$  is q-regular.

## From isolating Mahler equations to weighted automata

### Theorem (C., Yassawi, 2024)

Let  $q \ge 2$  be a natural number. There exists a universal q-automaton  $\mathcal{A}$ , such that any isolating q-Mahler equation P(x, y) over a commutative ring R with initial condition  $f_0$  provides weights for  $\mathcal{A}$ , so that the corresponding weighted automaton generates the solution f(x) of P(x, y) with  $f(0) = f_0$ .

▶ The universal q-automaton  $\mathcal{A}$  consists of a countable states set S and transitions in  $S \times \{0, 1, \dots, q-1\} \times S$ .

• Given an isolating q-Mahler equation  $P(x, y) = y - \sum_{i=1}^{d} \left( \sum_{j=0}^{h} \alpha_{i,j} x^{j} \right) \Phi^{i}(y)$ , we use its coefficients  $\alpha_{i,j}$  as weights, setting other edge weights to zero, so reducing  $\mathcal{A}$  to a weighted automaton.

## Example

For a 2-Mahler equation with height 3 and exponent 2

$$f(x) = A_1(x)f(x^2) + A_2(x)f(x^4)$$
  
where  $A_1(x) = \alpha_{1,0} + \alpha_{1,1}x + \alpha_{1,2}x^2 + \alpha_{1,3}x^3$   
and  $A_2(x) = \alpha_{2,0} + \alpha_{2,1}x + \alpha_{2,2}x^2 + \alpha_{2,3}x^3$ 



||

# Example

whilst for a 2-Mahler equation with height 3 and exponent 1

$$f(x) = (\alpha_{1,0} + \alpha_{1,1}x + \alpha_{1,2}x^2 + \alpha_{1,3}x^3)f(x^2)$$



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# From q-Mahler equations to Z-Mahler equations

## Theorem (Becker 1992, Dumas 1993)

Let  $q \ge 2$ , and let  $(u_n)_{n \ge 0}$  be a sequence over a commutative ring R.

- If  $(u_n)$  is q-regular, then it is the solution of a q-Mahler equation, and
- if  $(u_n)$  is a solution of an isolating q-Mahler equation, then it is q-regular.

## Theorem (C., Yassawi 2024)

Let  $(u_n)_{n \ge 0}$  be a sequence over a commutative ring R.

- ▶ If  $(u_n)$  is *q*-regular Z-regular, then it is the solution of a *q*-Mahler Z-Mahler equation,
- ▶ if  $(u_n)$  is a solution of an isolating  $\frac{q-Mahler}{Z-Mahler}$  Z-Mahler equation, then it is  $\frac{q-regular}{Z}$ -regular.

Our proof strategy was to emulate our proof in the case of q-numeration, i.e.,

▶ to define the linear Z-version of the map  $m \mapsto qm$ , and

▶ to define the appropriate concept of a Z-Mahler equation, in order to construct a weighted Z-automaton directly from an isolating Z-Mahler equation.

## The Zeckendorf analogue of $n \mapsto qn$

The map f(n) = qn can be written  $f(n) := [w0]_q$  where  $w = (n)_q$ . So, for  $(n)_Z = w$ , define  $\phi : \mathbb{N} \to \mathbb{N}$  as

Issue:  $\phi$  is not linear.

For example,

$$3 = \phi(2) = \phi(1+1) \neq 2\phi(1) = 4.$$

# Dealing with the nonlinearity of $\phi$

Recall  $\phi(n) := [(n)_Z 0]_Z$ . Define the linearity defect  $\delta$  by

$$\delta(m,n) := \phi(m+n) - \phi(m) - \phi(n).$$

A simple application of Binet's formula gives Lemma For natural numbers m, n, we have  $-1 \leq \delta(m, n) \leq 1$ .

In other words,  $\phi$  is almost linear.

We would like to track the linearity defect.

## Regularity of Z-expansions

Given a finite set  $C \subset \mathbb{Z}$ , consider

$$\mathcal{L}_C := \{ w \in C^* : [w]_Z = 0 \}$$

Example  $(C=\{0,1,-1\})$ 

The following words belong to  $\mathcal{L}_C$ :

55	34	21	13	8	5	3	2	1
0	1	-1	$^{-1}$	0	-1	1	1	0
1	-1	0	-1	0	-1	-1	0	0

### Theorem (Frougny)

For  $C \subset \mathbb{Z}$  finite,  $\mathcal{L}_C$  is a regular set of words.

### Corollary

There exists a deterministic automaton, which on input of  $(m)_Z$ and  $(n)_Z$ , outputs the linearity defect  $\delta(m-n,n)$  for  $m \ge n \ge 0$ .

## Going back to our strategy

We have defined the almost linear map  $m \mapsto \phi(m)$ .

We now define the Z-version of  $\Phi_q(\sum_n f_n x^n) = \sum_n f_n x^{qn}$ .

Define the Z-Mahler operator  $\Phi: R[\![x]\!] \to R[\![x]\!]$  as

$$\Phi\left(\sum_{n\geq 0} f_n x^n\right) := \sum_{n\geq 0} f_n x^{\phi(n)}.$$

The equation

$$P(x,y) = \sum_{i=0}^{d} A_i(x)\Phi^i(y) = 0$$

with  $A_i(x) \in R[x]$ , is a Z-Mahler equation.

If  $f \in R[x]$  satisfies  $\sum_{i=0}^{d} A_i(x) \Phi^i(f) = 0$ , then it is Z-Mahler.

Results

### Theorem (C., Yassawi 2024)

Let  $(u_n)_{n \ge 0}$  be a sequence over a commutative ring R.

- If  $(u_n)$  is Z-regular, then it is the solution of a Z-Mahler equation
- if  $(u_n)$  is a solution of an isolating Z-Mahler equation, then it is Z-regular.

Example  $(a_n = \# \text{ of Zeckendorf expansions of } n)$ 

$$f(x) = \sum_{n} a_n x^n = \prod_{n} (1 + x^{F_n})$$
 and  $f(x) = (1 + x)\Phi(f(x)).$ 



# Open questions

- ▶ Allouche and Shallit show that for  $a \in \mathbb{R}$ ,  $(a^n)_{n \ge 0}$  is *q*-regular if and only if a = 0 or *a* is a root of unity. Is there a similar result for Z-regular sequences?
- ▶ For R = C, Bell, Chyzak, Coons, & Dumas characterise q-regular series in terms of the q-Mahler equations they satisfy. Is there a similar characterisation for Zeckendorf numeration?
- Adamczewski-Bell and Shäfke-Singer show that a sequence which is both k- and  $\ell$ -Mahler over a field of characteristic zero, with k and  $\ell$  multiplicatively independent, must be rational. Which series are both k- and Z-Mahler?