

# Analogs of Markoff and Lagrange spectra on one-sided shift spaces

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- 1 Motivation: Diophantine approximation problems (simple case)
- 2 Markoff-Lagrange spectrum of shift spaces with cylinder order
- 3 Symmetric alphabets

# Diophantine approximation related to geometric progressions

## Fact

(1)  $\alpha \in \mathbb{R}$ ,  $\alpha > 1$ : fixed.

$\Rightarrow$  for almost every  $\xi \in \mathbb{R}$ ,  $(\xi\alpha^n)_{n \geq 0}$ : uniformly distributed modulo one (Weyl, 1916).

(2)  $b \in \mathbb{Z}$ ,  $b \geq 2$ : fixed.  $\xi \in \mathbb{R}$ ,  $\xi \geq 0$ .

Then  $(\xi b^n)_{n \geq 0}$ : uniformly distributed modulo one if and only if  $\xi$  is normal in base  $b$ .

However, for each individual  $\xi$ , it is generally difficult to investigate the distribution  $\xi b^n$  modulo one.

## Target ( $b \in \mathbb{Z}$ , $b \geq 2$ )

$\limsup_{n \rightarrow \infty} \|\xi b^n\|$ , where  $\|x\| = \min\{|x - m| \mid m \in \mathbb{Z}\} \in [0, 1/2]$ .

Example:  $\|5.9\| = |5.9 - 6| = |-0.1| = 0.1$ .

$\mathcal{L}(b) := \{\limsup_{n \rightarrow \infty} \|\xi b^n\| \mid \xi \in \mathbb{R}\} \subset [0, 1/2]$ .

# Multiplicative Lagrange spectrum

Target ( $b \in \mathbb{Z}, b \geq 2$ )

$\limsup_{n \rightarrow \infty} \|\xi b^n\|$ , where  $\|x\| = \min\{|x - m| \mid m \in \mathbb{Z}\} \in [0, 1/2]$ .  
 $\mathcal{L}(b) = \{\limsup_{n \rightarrow \infty} \|\xi b^n\| \mid \xi \in \mathbb{R}\} \subset [0, 1/2]$ .

Theorem (Main idea was obtained by Dubickas 2006)

$(\exists C(b) \in \mathcal{L}(b))$  satisfying the following:

- ①  $(0, C(b)) \cap \mathcal{L}(b)$  is an infinite discrete set denoted as  
 $\{d^{(b)}(0) < d^{(b)}(1) < d^{(b)}(2) < \dots\}$ , where  $\lim_{m \rightarrow \infty} d^{(b)}(m) = C(b)$ .  
In particular,  $C(b)$  is the minimal limit point of  $\mathcal{L}(b)$ .
- ②  $(\forall m \geq 0) d^{(b)}(m) \in \mathbb{Q}$ ,  $C(b)$  is a transcendental number.

We call  $[0, C(b)] \cap \mathcal{L}(b)$  the **discrete part of  $\mathcal{L}(b)$** .

# Discrete part of $\mathcal{L}(b)$ and substitution ( $b \in \mathbb{Z}, b \geq 2$ )

$\mathcal{L}(b) = \{\limsup_{n \rightarrow \infty} \|\xi b^n\| \mid \xi \in \mathbb{R}\} \subset [0, 1/2]$ .  $C(b)$ : min. lim. p.t. of  $\mathcal{L}(b)$ .  
 $(0, C(b)) \cap \mathcal{L}(b) = \{d^{(b)}(0) < d^{(b)}(1) < d^{(b)}(2) < \dots\}$

## Substitution and additional rule

(1)  $\tau : \{0, 1\}^* \rightarrow \{0, 1\}^*$ .  $\tau(0) = 1, \tau(1) = 100$ .

$A_0 := 0, A_{m+1} := \tau(A_m)$ .  $A_1 = 1, A_2 = 100, A_3 = 10011, \dots$

$A_\infty := \lim_{m \rightarrow \infty} A_m = 10011100100\dots$ : fixed point of  $\tau$ .

(2) For  $x = x_0x_1\dots \in \{0, 1\}^\infty$ , put  $\Phi(x) := 10^{x_0}\bar{1}0^{x_1}10^{x_2}\bar{1}0^{x_3}\dots$ , where  $\bar{1} = -1$ .

If  $x_n = 1$ , then insert 0.  $v_m := \Phi(A_m^\infty) = \Phi(A_m A_m A_m \dots)$ ,  $v_\infty := \Phi(A_\infty)$ .

Example:  $A_2^\infty = 100100100\dots v_2 = 10\bar{1}1\bar{1}01\bar{1}10\bar{1}11\dots$

## "base- $b$ " expansion of $d^{(b)}(m)$ and $C(b)$

$d^{(b)}(m) = (.v_m)_b, C(b) = (.v_\infty)_b$ . For instance,

$d^{(b)}(2) = (.v_2)_b = (.10\bar{1}1\bar{1}01\bar{1}10\bar{1}11\dots)_b = b^{-1} - b^{-3} + b^{-4} - b^{-5} + b^{-7} - b^{-8} + \dots$ .

# Examples of $\limsup_{n \rightarrow \infty} \|\xi b^n\|$ ( $b \in \mathbb{Z}, b \geq 2$ )

For simplicity, we assume that  $b \geq 3$

$$\xi = \sum_{m=1}^{\infty} b^{-m^2} = (.100100001000 \dots)_b = \|\xi\|.$$

$$\xi b = (1.00100001000 \dots)_b, \|\xi b\| = (.00100001000 \dots)_b,$$

$$\xi b^2 = (10.0100001000 \dots)_b, \|\xi b^2\| = (.0100001000 \dots)_b \dots$$

Key obs.: Distribution of  $\|\xi b^n\|$  ( $n = 0, 1, \dots$ )  $\Rightarrow$  shifts of the word 100100001000 ...

$$\limsup_{n \geq 0} \|\xi b^n\| = (.10000 \dots)_b = 1/b \text{ because } \lim_{m \rightarrow \infty} ((m+1)^2 - m^2) = \infty.$$

Interpretation in terms of seq.:  $w := 100100001000 \dots$  (seq. cor. to base- $b$  exp. of  $\xi$ ).

Then limit sup word (**limsup word**) of  $w$  w.r.t. **lexicographical order**

is 1000... =  $10^\infty$ . (cf.  $\sup_{n \geq 0} \|\xi b^n\| = \|\xi\| = (.100100001000 \dots)_b$ .)

## Our goal

Limsup words w.r.t. one-sided spaces with general order (for further application).

In particular, research the minimal limit point in terms of substitutions.

# Cylinder order on $\{0, 1\}^\infty$

$\mathcal{A} := \{0, 1\}$ ,  $\mathcal{A}^\infty := \{\mathbf{a} = a_1 a_2 \cdots \mid a_1, a_2, \dots \in \mathcal{A}\}$ .

$\mathbf{a} = a_1 a_2 \cdots, \mathbf{b} = b_1 b_2 \cdots \in \mathcal{A}^\infty \Rightarrow d(\mathbf{a}, \mathbf{b}) := 2^{-k}$ , where  $k = \min\{n \geq 1 \mid a_n \neq b_n\}$ .

$a_1, \dots, a_m \in \mathcal{A} \Rightarrow [\mathbf{a}_1 \cdots \mathbf{a}_m] := \{a'_1 a'_2 \cdots \in \mathcal{A}^\infty \mid a'_1 \cdots a'_k = a_1 \dots a_k\}$ : cylinder set.

Cylinder order on  $\mathcal{A}^\infty$  (or  $\{\bar{1}, 0, 1\}^\infty$  in the next section)

$\leq$ : total order on  $\mathcal{A}^\infty$ .

$\leq$  is cylinder order if the following holds:

For any  $\mathbf{a}, \mathbf{b}, \mathbf{c} \in \mathcal{A}^\infty$ ,  $\mathbf{a} \leq \mathbf{b} \leq \mathbf{c}$  implies  $d(\mathbf{a}, \mathbf{b}) \leq d(\mathbf{a}, \mathbf{c})$ .

Equivalent definition of cylinder order:  $k \geq 1$ ,  $a_1, \dots, a_k, b_1, \dots, b_k \in \mathcal{A}$ .

We define  $[a_1 \cdots a_k] < [b_1 \cdots b_k]$  by  $\mathbf{a} < \mathbf{b}$  for any  $\mathbf{a} \in [a_1 \cdots a_k]$  and  $\mathbf{b} \in [b_1 \cdots b_k]$ .

Total order on  $\mathcal{A}^\infty$  is cylinder order if and only if, for any  $k \geq 1$ ,

$<$  is a total order on  $\{[a_1 \cdots a_k] \mid a_1, \dots, a_k \in \mathcal{A}\}$  (the set of cyl. set of length  $k$ ).

Example: lexicographic order (next page).

# limsup words

$\mathcal{A} = \{0, 1\}$ ,  $\mathcal{A}^\infty = \{\mathbf{a} = a_1 a_2 \cdots \mid a_1, a_2, \dots \in \mathcal{A}\}$ .  $[a_1 \cdots a_k]$  : cylinder set.

$\leq$ : cylinder order  $\Leftrightarrow (\forall k \geq 1) <$  is a total order on  $\{[a_1 \cdots a_k] \mid a_1, \dots, a_k \in \mathcal{A}\}$ .

Examples:  $\mathbf{a} = a_1 a_2 \cdots$ ,  $\mathbf{b} = b_1 b_2 \cdots \in \mathcal{A}^\infty$  with  $\mathbf{a} \neq \mathbf{b}$ .  $m := \min\{n \geq 1 \mid a_n \neq b_n\}$ .

(1) Lexicographic order  $\leq_{lex}$ :  $\mathbf{a} <_{lex} \mathbf{b} \Leftrightarrow a_m < b_m$ .

$\Rightarrow [0] <_{lex} [1]$ ,  $[00] <_{lex} [01] <_{lex} [10] <_{lex} [11]$ .

(2) Alternating lexicographic order  $\leq_{alt}$ :  $\mathbf{a} <_{alt} \mathbf{b} \Leftrightarrow \begin{cases} a_m < b_m & m \text{ is odd}, \\ a_m > b_m & m \text{ is even}. \end{cases}$

$\Rightarrow [0] <_{alt} [1]$ ,  $[01] <_{alt} [00] <_{alt} [11] <_{alt} [10]$ ,  $[100] <_{alt} [101]$ .

## limsup words, sup words

$\mathbf{a} = a_1 a_2 \cdots \in \mathcal{A}^\infty$ .

$\ell_{\leq}(\mathbf{a}) := \limsup_{n \rightarrow \infty} a_n a_{n+1} a_{n+2} \cdots$ ,  $s_{\leq}(\mathbf{a}) := \sup_{n \geq 1} a_n a_{n+1} a_{n+2} \cdots$ .

Example:  $\mathbf{a} = 10010000100000010 \cdots$  ( $a_n = 1$  iff.  $n$  is a square number  $c^2$ .)

$\ell_{\leq_{lex}}(\mathbf{a}) = 1000000 \cdots = 10^\infty$ ,  $s_{\leq_{lex}}(\mathbf{a}) = \mathbf{a}$ .

# Analogs of Markoff-Lagrange spectrum

$\mathcal{A} = \{0, 1\}$ ,  $\mathcal{A}^\infty = \{\mathbf{a} = a_1 a_2 \cdots \mid a_1, a_2, \dots \in \mathcal{A}\}$ .  $[a_1 \cdots a_k]$  : cylinder set.

$\leq$ : cylinder order  $\Leftrightarrow (\forall k \geq 1) <$  is a total order on  $\{[a_1 \cdots a_k] \mid a_1, \dots, a_k \in \mathcal{A}\}$ .

$\mathbf{a} = a_1 a_2 \cdots \in \mathcal{A}^\infty$ .

$\ell_{\leq}(\mathbf{a}) = \limsup_{n \rightarrow \infty} a_n a_{n+1} a_{n+2} \cdots$ ,  $s_{\leq}(\mathbf{a}) = \sup_{n \geq 1} a_n a_{n+1} a_{n+2} \cdots$ .

## Definition

$\mathcal{L}_{\leq} := \{\ell_{\leq}(\mathbf{a}) \mid \mathbf{a} \in \mathcal{A}^\infty\}$ : Lagrange spectrum w.r.t.  $\leq$ .

$\mathcal{M}_{\leq} := \{s_{\leq}(\mathbf{a}) \mid \mathbf{a} \in \mathcal{A}^\infty\}$ : Markoff spectrum w.r.t.  $\leq$ .

## Theorem( $\leq$ : cylinder order)

$\mathcal{L}_{\leq} = \mathcal{M}_{\leq} = \text{closure}\{s_{\leq}(\mathbf{a}) \mid \mathbf{a} \in \mathcal{A}^\infty \text{ is purely periodic.}\}$ .

Recall that  $\mathcal{A}^\infty$  is a metric space with metric  $d(\mathbf{a}, \mathbf{b}) = 2^{\min\{n \geq 1 \mid a_n \neq b_n\}}$ .

Next target: minimal limit point  $\mathbf{m}_{\leq}$  of  $\mathcal{L}_{\leq} = \mathcal{M}_{\leq}$ .

# Notation for substitution

$$\mathcal{A} = \{0, 1\}, \mathcal{A}^\infty = \{\mathbf{a} = a_1 a_2 \cdots \mid a_1, a_2, \dots \in \mathcal{A}\}.$$

$\leq$ : cylinder order  $\Leftrightarrow (\forall k \geq 1) <$  is a total order on  $\{[a_1 \cdots a_k] \mid a_1, \dots, a_k \in \mathcal{A}\}$ .

$$\mathcal{L}_\leq = \mathcal{M}_\leq = \{\sup_{n \geq 1} a_n a_{n+1} \cdots \mid a_1 a_2 \cdots \in \mathcal{A}^\infty\}.$$

## Definition

- (1)  $\mathbf{m}_\leq$ : minimal limit point of  $\mathcal{L}_\leq = \mathcal{M}_\leq$ .
- (2)  $0 \leq j < k \Rightarrow \tau_{j,k} : \{0, 1\}^* \rightarrow \{0, 1\}^*$ : substitution with  $\tau_{j,k}(0) = 10^j, \tau_{j,k}(1) = 10^k$ .  
Note that  $\tau$  ( $\tau(0) = 1, \tau(1) = 100$ ) in the result of Dubickas 2006 is  $\tau_{0,2}$ .
- (3)  $\mathcal{S} = \{\tau_{j,k} \mid 0 \leq j < k\}$ .
- (4)  $\sigma = (\sigma_n)_{n \geq 1} \in \mathcal{S}^\infty$ .  $\sigma_{[1,n]} := \sigma_1 \circ \sigma_2 \circ \cdots \circ \sigma_n$ ,  $\sigma := \lim_{n \rightarrow \infty} \sigma_{[1,n]}$ .

Next page: We describe  $\{\mathbf{a} \in \mathcal{M}_\leq \mid \mathbf{a} \leq \mathbf{m}_\leq\}$ : Analog of discrete part.

We may assume that  $[0] < [1]$ .

## Discrete part

$\mathcal{A} = \{0, 1\}$ ,  $\mathcal{A}^\infty = \{\mathbf{a} = a_1 a_2 \cdots \mid a_1, a_2, \dots \in \mathcal{A}\}$ .  $\leq$ : cylinder order with  $[0] < [1]$ .

$\mathbf{m}_\leq$ : min. lim. p.t. of  $\mathcal{L}_\leq = \mathcal{M}_\leq = \{\sup_{n \geq 1} a_n a_{n+1} \cdots \mid a_1 a_2 \cdots \in \mathcal{A}^\infty\}$ .

$\mathcal{S} = \{\tau_{j,k} \mid 0 \leq j < k\}$ ,  $\tau_{j,k}(0) = 10^j$ ,  $\tau_{j,k}(1) = 10^k$ .

$\sigma = (\sigma_n)_{n \geq 1} \in \mathcal{S}^\infty$ .  $\sigma_{[1,n]} := \sigma_1 \circ \sigma_2 \circ \cdots \circ \sigma_n$ ,  $\sigma := \lim_{n \rightarrow \infty} \sigma_{[1,n]}$ .

Theorem (there is an algorithm to determine  $\sigma_1, \sigma_2, \dots$ )

Exactly one of the following holds:

(1) ( $\exists \sigma = (\sigma_n)_{n \geq 1} \in \mathcal{S}^\infty$ ) s.t.  $\mathbf{m}_\leq = \sigma(1^\infty)$ . Moreover, the following holds:

$\{\mathbf{a} \in \mathcal{M}_\leq \mid \mathbf{a} < \mathbf{m}_\leq\} = \{(\sigma_{[1,n]}(0))^\infty \mid n \geq 0\}$ , where if  $n = 0$ ,  $(\sigma_{[1,n]}(0))^\infty = 0^\infty$ .

(2) There exists  $h \geq 0$  and  $\sigma_1, \dots, \sigma_h \in \mathcal{S}$  such that  $\mathbf{m}_\leq = \sigma_{[1,h]}(10^\infty)$ , where  $\sigma_{[1,h]}(10^\infty) = 10^\infty$  when  $h = 0$ . Moreover,  $\{\mathbf{a} \in \mathcal{M}_\leq \mid \mathbf{a} < \mathbf{m}_\leq\}$  is as follows:

$$\{(\sigma_{[1,n]}(0))^\infty \mid 0 \leq n \leq h\} \cup \{\sigma_{[1,h]}((10^j)^\infty) \mid j \geq 0, \sigma_{[1,h]}((10^j)^\infty) < \mathbf{m}_\leq\}.$$

## Discrete part (Example)

$\mathcal{A} = \{0, 1\}$ ,  $\mathcal{A}^\infty = \{\mathbf{a} = a_1 a_2 \cdots \mid a_1, a_2, \dots \in \mathcal{A}\}$ .  $\leq$ : cylinder order with  $[0] < [1]$ .  
 $\mathbf{m}_\leq$ : min. lim. p.t. of  $\mathcal{L}_\leq = \mathcal{M}_\leq$ .  $\mathcal{S} = \{\tau_{j,k} \mid 0 \leq j < k\}$ ,  $\tau_{j,k}(0) = 10^j$ ,  $\tau_{j,k}(1) = 10^k$ .  
 $\boldsymbol{\sigma} = (\sigma_n)_{n \geq 1} \in \mathcal{S}^\infty$ .  $\sigma_{[1,n]} = \sigma_1 \circ \sigma_2 \circ \cdots \circ \sigma_n$ ,  $\boldsymbol{\sigma} = \lim_{n \rightarrow \infty} \sigma_{[1,n]}$ .

Theorem (The second statement is omitted.)

$(\exists \boldsymbol{\sigma} = (\sigma_n)_{n \geq 1} \in \mathcal{S}^\infty)$  s.t.  $\mathbf{m}_\leq = \boldsymbol{\sigma}(1^\infty)$ . Moreover, the following holds:  
 $\{\mathbf{a} \in \mathcal{M}_\leq \mid \mathbf{a} < \mathbf{m}_\leq\} = \{(\sigma_{[1,n]}(0))^\infty \mid n \geq 0\}$ , where if  $n = 0$ ,  $(\sigma_{[1,n]}(0))^\infty = 0^\infty$ .

Alternating lexicographic order  $\leq_{alt}$

$\mathbf{a} = a_1 a_2 \cdots$ ,  $\mathbf{b} = b_1 b_2 \cdots \in \mathcal{A}^\infty$  with  $\mathbf{a} \neq \mathbf{b}$ .  $m := \min\{n \geq 1 \mid a_n \neq b_n\}$ .

$\mathbf{a} <_{alt} \mathbf{b} \Leftrightarrow a_m < b_m (m \text{ is odd}), a_m > b_m (m \text{ is even})$ .

$(\forall n \geq 1) \sigma_n = \tau_{0,2}$ .  $\sigma_{[1,n]} = \tau_{0,2}^n$ ,  $\boldsymbol{\sigma}(1^\infty) = \lim_{n \rightarrow \infty} \tau_{0,2}^n(1^\infty) = \lim_{n \rightarrow \infty} \tau_{0,2}^n(1)$ .

# Unimodal order (If we have no time, then we omit this part)

$\mathcal{A} = \{0, 1\}$ ,  $\mathcal{A}^\infty = \{\mathbf{a} = a_1 a_2 \cdots \mid a_1, a_2, \dots \in \mathcal{A}\}$ .  $\leq$ : cylinder order with  $[0] < [1]$ .  
 $\mathbf{m}_{\leq}$ : min. lim. p.t. of  $\mathcal{L}_{\leq} = \mathcal{M}_{\leq}$ .  $\mathcal{S} = \{\tau_{j,k} \mid 0 \leq j < k\}$ ,  $\tau_{j,k}(0) = 10^j$ ,  $\tau_{j,k}(1) = 10^k$ .  
 $\boldsymbol{\sigma} = (\sigma_n)_{n \geq 1} \in \mathcal{S}^\infty$ .  $\sigma_{[1,n]} = \sigma_1 \circ \sigma_2 \circ \cdots \circ \sigma_n$ ,  $\boldsymbol{\sigma} = \lim_{n \rightarrow \infty} \sigma_{[1,n]}$ .

Unimodal order  $\leq_{uni}$  (related to the dynamics of unimodal maps, e.g. tent map.)

$\mathbf{a} = a_1 a_2 \cdots, \mathbf{b} = b_1 b_2 \cdots \in \mathcal{A}^\infty$  with  $\mathbf{a} \neq \mathbf{b}$ .  $m := \min\{n \geq 1 \mid a_n \neq b_n\}$ .  
 $N := \text{Card}\{1 \leq n \leq m-1 \mid a_n = 1\}$  (if  $m=1$ , then  $N:=0$ ).  
 $\mathbf{a} <_{uni} \mathbf{b} \Leftrightarrow a_m < b_m$  ( $N$  is even),  $a_m > b_m$  ( $N$  is odd).

Minimal limit point  $\mathbf{m}_{\leq_{uni}}$

$\sigma_1 := \tau_{0,1}$ ,  $(\forall n \geq 2) \sigma_n := \tau_{0,2}$ .  
 $\sigma_{[1,1]} = \tau_{0,1}$ ,  $(\forall n \geq 2) \sigma_{[1,n]} = \tau_{0,1} \circ \tau_{0,2}^{n-1}$ ,  $\boldsymbol{\sigma}(1^\infty) = \lim_{n \rightarrow \infty} \tau_{0,1} \circ \tau_{0,2}^{n-1}(1)$ .

# Consistent cylinder order on $\{\bar{1}, 0, 1\}^\infty$

$\mathcal{B} := \{\bar{1}, 0, 1\}$ ,  $\mathcal{B}^\infty := \{\mathbf{a} = a_1 a_2 \cdots \mid a_1, a_2, \dots \in \mathcal{B}\}$ .

Model (motivation):  $\limsup_{n \rightarrow \infty} \left| \sum_{m=1}^{\infty} a_{n+m} \beta^{-m} \right|$  ( $\beta \geq 3$ ).

More generally,  $(e_m)_{m=1,2,\dots} \in \{\bar{1}, 1\}^\infty$ : fixed sequence of signature.

$\Rightarrow \limsup_{n \rightarrow \infty} \left| \sum_{m=1}^{\infty} e_m \cdot a_{n+m} \beta^{-m} \right|$  ( $\beta \geq 3$ ).

## Absolute values of sequences

$$\mathbf{a} \in \mathcal{B}^\infty \Rightarrow \text{abs}(\mathbf{a}) := \begin{cases} \mathbf{a} & \text{if } \mathbf{a} \geq_{lex} 0^\infty, \\ -\mathbf{a} & \text{if } \mathbf{a} \leq_{lex} 0^\infty. \end{cases}$$

## Definition

$\preceq$ : cylinder order on  $\mathcal{B}^\infty$ .  $\preceq$  is **consistent** if

$(\forall w \in \mathcal{B}^*) "[w\bar{1}] \prec [w0] \prec [w1]" \text{ or } "[w1] \prec [w0] \prec [w\bar{1}]".$

# Markoff-Lagrange spectrum on $\{\bar{1}, 0, 1\}^\infty$

$\mathcal{B} = \{\bar{1}, 0, 1\}$ ,  $\mathcal{B}^\infty = \{\mathbf{a} = a_1 a_2 \cdots \mid a_1, a_2, \dots \in \mathcal{B}\}$ .  $\preceq$ : **consistent** cyl. order on  $\mathcal{B}^\infty$ :  
 $(\forall w \in \mathcal{B}^*) "[w\bar{1}] \prec [w0] \prec [w1]" \text{ or } "[w1] \prec [w0] \prec [w\bar{1}]".$

$\mathbf{a} \in \mathcal{B}^\infty : \text{abs}(\mathbf{a}) = \mathbf{a}$  if  $\mathbf{a} \geq_{lex} 0^\infty$ ,  $\text{abs}(\mathbf{a}) = -\mathbf{a}$  if  $\mathbf{a} \leq_{lex} 0^\infty$ .

Definition ( $\preceq$ : cylinder order)

(1)  $\mathbf{a} = a_1 a_2 \cdots \in \mathcal{B}^\infty$ .

$\ell_{\preceq}^{\text{abs}}(\mathbf{a}) := \limsup_{n \rightarrow \infty} \text{abs}(a_n a_{n+1} a_{n+2} \cdots)$ ,  $s_{\preceq}^{\text{abs}}(\mathbf{a}) := \sup_{n \geq 1} \text{abs}(a_n a_{n+1} a_{n+2} \cdots)$ .

(2)  $\mathcal{L}_{\preceq}^{\text{abs}} := \{\ell_{\preceq}^{\text{abs}}(\mathbf{a}) \mid \mathbf{a} \in \mathcal{B}^\infty\}$ .  $\mathcal{M}_{\preceq}^{\text{abs}} := \{s_{\preceq}^{\text{abs}}(\mathbf{a}) \mid \mathbf{a} \in \mathcal{B}^\infty\}$ .

Theorem ( $\preceq$ : cylinder order, not necessary consistent)

$\mathcal{L}_{\preceq}^{\text{abs}} = \mathcal{M}_{\preceq}^{\text{abs}} = \text{closure}\{s_{\preceq}^{\text{abs}}(\mathbf{a}) \mid \mathbf{a} \in \mathcal{B}^\infty \text{ is purely periodic}\}$ ,

where  $\mathcal{B}^\infty$  is a metric space with  $d(\mathbf{a}, \mathbf{b}) = 2^{\min\{n \geq 1 \mid a_n \neq b_n\}}$ .

Next target: minimal limit point  $\mu_{\preceq}$  of  $\mathcal{L}_{\preceq}^{\text{abs}} = \mathcal{M}_{\preceq}^{\text{abs}}$ .

## Discrete part

$\mathcal{B} = \{\bar{1}, 0, 1\}$ ,  $\mathcal{B}^\infty = \{a = a_1 a_2 \cdots \mid a_1, a_2, \dots \in \mathcal{B}\}$ .  $\preceq$ : consistent cyl. order on  $\mathcal{B}^\infty$ :  
 $(\forall w \in \mathcal{B}^*) "[w\bar{1}] \prec [w0] \prec [w1]" \text{ or } "[w1] \prec [w0] \prec [w\bar{1}]".$

$a \in \mathcal{B}^\infty$ :  $\text{abs}(a) = a$  if  $a \geq_{lex} 0^\infty$ ,  $\text{abs}(a) = -a$  if  $a \leq_{lex} 0^\infty$ .

$\mu_\preceq$ : min. lim. p.t. of  $\mathcal{L}_\preceq^{\text{abs}} = \{\limsup_{n \rightarrow \infty} \text{abs}(a_n a_{n+1} a_{n+2} \cdots) \mid a \in \mathcal{B}^\infty\}$ .

Recall for  $x = x_0 x_1 \dots \in \{0, 1\}^\infty$  that  $\Phi(x) = 10^{x_0} \bar{1} 0^{x_1} 10^{x_2} \bar{1} 0^{x_3} \dots$

For simplicity, we only consider the case where  $[1\bar{1}] \prec [10]$  and  $[10\bar{1}] \prec [100]$  hold.

Theorem ( $\preceq$ : with assumption above.)

There exists a cylinder order  $\leq$  on  $\mathcal{A}^\infty = \{0, 1\}^\infty$  such that  $\mu_\preceq = \Phi(m_\leq)$ ,  
where  $m_\leq$  is the min. lim. p.t. of  $\mathcal{L}_\leq$ . Moreover, the following holds:

$$\{a \in \mathcal{L}_\preceq^{\text{abs}} \mid a \prec \mu_\preceq\} = \{0^\infty\} \cup \{\Phi(b) \mid b \in \mathcal{L}_\leq, b < m_\leq\}.$$

Note: there exists an algorithm to determine  $\leq$  by using  $\preceq$ .

Example:  $\preceq_{lex}$ : lexicographic order on  $\mathcal{B}^\infty$ : next page.

## Example

$$\mathcal{B} = \{\bar{1}, 0, 1\}, \mathcal{B}^\infty = \{a = a_1 a_2 \cdots \mid a_1, a_2, \dots \in \mathcal{B}\}.$$

$\preceq$ : consistent cyl. order on  $\mathcal{B}^\infty$ .  $\mu_{\preceq}$ : min. lim. p.t. of  $\mathcal{L}_{\preceq}^{\text{abs}}$ .

$$\mathbf{x} = x_0 x_1 \dots \in \{0, 1\}^\infty \Rightarrow \Phi(\mathbf{x}) = 10^{x_0} \bar{1} 0^{x_1} 1 0^{x_2} \bar{1} 0^{x_3} \dots$$

Theorem (assumption:  $\preceq$  satisfies  $[1\bar{1}] \prec [10]$ ,  $[10\bar{1}] \prec [100]$ )

There exists a cylinder order  $\leq$  on  $\mathcal{A}^\infty = \{0, 1\}^\infty$  such that  $\mu_{\preceq} = \Phi(\mathbf{m}_{\leq})$ ,

where  $\mathbf{m}_{\leq}$  is the min. lim. p.t. of  $\mathcal{L}_{\leq}$ . Moreover, the following holds:

$$\{a \in \mathcal{L}_{\preceq}^{\text{abs}} \mid a \prec \mu_{\preceq}\} = \{0^\infty\} \cup \{\Phi(b) \mid b \in \mathcal{L}_{\leq}, b < \mathbf{m}_{\leq}\}.$$

Example:  $\preceq_{lex}$ : lexicographic order on  $\mathcal{B}^\infty \Rightarrow \leq_{alt}$ : alternative lex. order.

$\mu_{\preceq_{lex}} = \Phi(\mathbf{m}_{\leq_{alt}}) = \Phi(\lim_{n \rightarrow \infty} \tau_{0,2}^n(1))$  (Result by Dubickas 2006).

We also have other examples (alt. lex. order on  $\mathcal{B}^\infty$  and bimodal order on  $\mathcal{B}^\infty$ ).